# Enhancing a Greek Language Stemmer 

Efficiency and Accuracy Improvements

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Stemming algorithms are used in the field of Information Retrieval in order to improve precision and recall. Although for Greek there are three stemmers published, only one of them is freely available. In this thesis, we use stemmer performance metrics for evaluating the existing algorithm and we improved its accuracy and completeness. These improvements were achieved by providing an alternative implementation in PHP which offers more syntactical rules and exceptions. Finally, the two algorithms are tested and their statistics metrics are compared.

Key words and terms: stemming algorithm, Greek, algorithmic efficiency, PHP

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## 1. Introduction

Search engines play a critical role in people's life nowadays. Google reported back in 2006 that it receives more than 100 million queries from US based hosts daily [Witten et al., 2007]. For end users, they are perhaps the only way to find the information they seek and to navigate to the appropriate web pages. In addition, the users of search engines tend to navigate through the top results that the search engines return to them. It is quite common that end users know which piece and what type of information they seek but they are quite often unable to construct a proper query that will fully describe their request. As a result, malformed and badly structured queries tend to return fewer and more irrelevant results than well/structured queries. For example, someone seeking information for "World War II" may form a query as "World Wars". Since "Wars" is not a part of "War", documents containing the reference "War" will not be returned and, thus, the number of relevant results will be reduced. One of the solutions proposed to increase the ratio of relevant documents to total documents retrieved, also known as recall, is stemming.

### 1.1 Stemming

Stemming is the process of reducing words to their stem, base or root form, [Lovins, 1968] as shown in Table 1.

Table 1: Stemming Examples

| in English |  |
| :--- | :--- |
| Original word |  |
| Seas | sea |
| Wars | war |
| determination | determin- |
| Developed | develop |
|  | in Greek |
| Avaлvと́ $\omega$ | $\alpha v \alpha \pi v \varepsilon ́-~$ |
| E $\lambda \lambda \eta \nu ı$ кós | $\varepsilon \lambda \lambda \eta v i \kappa-$ |

This is a simple but effective operation used in the fields of information extraction and natural languages [Carlberger et al., 2001]. Stemming can be utilized when storing
information about a web page in a search engine's database or for query expansion $[\mathrm{Xu}$ \& Croft, 1998]. In that case, the user's query is evaluated and reformulated. The search engine may reduce the word to its stem and thus return more results to the user. An evaluation research in 1981 showed that stemming improved search precision [Brants, 2003].

### 1.2 Definitions of Key Terms and Concepts

As stemming is also a linguistic process, any discussion of it, even in the field of IR, assumes the knowledge of some basic linguistic terms. The ones used in this thesis are described in Table 2.

Table 2: Definitions of key terms in the stemming process
Stem: a base part of a word that may have or may not have semantic meaning and no affixes (see below)

Affix: a small linguistic unit with semantic meaning that is attached to the beginning or the end of a stem to form a word

Prefix: an affix that is added to the beginning of a word or stem
Suffix: an affix that is added to the end of a word or stem
Stop Word: a term that appears so frequently in documents that it does not help searches [van Rijsbergen, 1979]. For example:

I, a, to, any, where, you
Inflection: the modification or marking of a word to reflect semantic and grammatical information like gender, tense, number, case or person.

For example:
to help - > helps, helping, helped (reflects gender and tense)
Derivation : the modification of a word that transforms it from one syntactical category (verb, noun, adverb, adjective) to another.

For example:
to hope - > hopefully, hopeless, hopes (transformation to adverb, adjective and noun/verb)

Conflation class: a group of words that share the same semantic meaning [Paice, 1996]. For example:
group, grouping, teams
Compounding: the creation of new words by combination of two or more different words into a single form. For example:
solar-powered, breastfeeding, bitter-sweet, antidisestablishmentarianism
Morphological Variants: two or more words that are related due to inflation, derivation of compounding.

For example:
dark, darkening, darks, darkroom

### 1.3 Inflectional versus Derivational Variants

Inflectional and derivational are the two categories of morphological variants that stemmers mainly deal with. Any given word can have inflectional and derivational morphological variants. Inflectional morphological variants share the same basic meaning and belong to the same part of speech. For instance, the changes may affect the word's case, number, tense and gender. In contrast derivational morphological variants can belong to a different part of speech and thus, may mean something completely different from their stem since an adjective can be derived from a noun, or a noun from a verb among others. Many of the algorithms so far developed do not treat derivational suffixes or handle them partially. According to Paice [1994], affixes may contain important information about the meaning of the word and so it is advisable not to discard it during the stemming process. Stemming "antidote" to "dote " creates a word that belongs to a different conflation class since the two words deal with different concepts.

Paice refers to the English language. However, the same phenomenon is noticed also in Greek. Affixes that are added usually alter the meaning of the word completely compared to the initial meaning of the stem. Some examples of this type of derivation are " $\mu$ ó $\varphi \varphi \sigma \eta "$ " (education) to " $\pi \alpha \rho \alpha \mu о ́ \rho \varphi \omega \sigma \eta " ~(d i s f i g u r e m e n t, ~ d e f o r m a t i o n) ~ a n d ~$ " $\delta$ окци́ऽ $\omega$ " (I try) to " $\alpha \pi о \delta о к ц \alpha ́ \zeta \omega " ~(I ~ b o o, ~ I ~ a b j u r e) . ~ A ~ s t e m m e r ~ t h a t ~ i n c l u d e s ~$ derivational rules can be helpful for language research, but may be inappropriate as a query expansion tool, since it can supply the search engine with variants of the original
words that have a very different meaning. This fact can increase the number of matched documents but the semantic accuracy to the original word will be low.

### 1.4 Stemming Techniques - Advantages and Disadvantages

In addition to whether or not a stemmer treats prefixes, they are also categorized as (i) dictionary based (ii) based on algorithms or (iii) a hybrid version of both [Ntais, 2006]. Dictionary based stemmers use ready made dictionaries and match a word with its stem from a list. The main drawback of such stemmers is that dictionary maintenance is required and that these stemmers can not scale to handle unlimited words.

In contrast, algorithm based stemmers have been the focus of research with the algorithms of Lovins and Porter being the most representative. The former [Lovins, 1968] precedes the latter by 12 years and it was the first stemming algorithm ever published. It uses an extensive list of 294 endings, 35 transformation rules and 29 conditions. The algorithm is executed in two basic steps. In the first step the longest ending is matched and removed and in the second step the algorithm checks whether one of the 35 transformation rules should be applied.

Although Lovins' algorithm was the first stemming algorithm, Porter's is considered one of the most influential [Krovetz, 1993; Ntais, 2006; Frakes, 2003]. It was initially written for the English language and later ported to other European languages like Italian, Spanish, French and Portuguese. In contrast with Lovins' algorithm, it iteratively applies a set of rules and removes suffixes until no rules apply. The execution is completed into five distinct steps [Porter, 1980] and is considered very aggressive [Krovetz, 1993; Xu and Croft 1998]. Porter's work became influential and many implementations were written and made available by others. Unfortunately, these implementations contained errors. In order to deal with this problem, Porter released a framework with which stemming algorithms can be implemented using a string handling programming language called Snowball.

### 1.5 Greek Stemmers

Three stemmers have been developed for Greek. The first two are the TZK algorithm by Kalamboukis and Nikolaidis in 1995 and the Automated Morphological Processor (AMP) by Tambouratzis and Carayanis in 2001. These two stemmers have an acceptable accuracy of 90 to $95 \%$ but for their development some constraints were implied. The

AMP algorithm assumes that each word just consists of a stem part and an ending part and thus excludes all compound words. One the other hand, the TZK algorithm can only manipulate 65 suffixes although there are more than 166 suffixes in the Greek language. Furthermore, to our knowledge, neither of these algorithms has a freely available implementation.

The latest stemming algorithm developed for the Greek language is by George Ntais, in 2006. This algorithm follows the structure of the Porter algorithm and has a free implementation available on the web [Ntais, 2008]. The author has provided a web interface where users can make simple queries by posting a single word in Greek and have the word's stem returned. The interface is simple and uses Javascript for the implementation of the algorithm. According to the author, the algorithm can handle 158 suffixes of the Greek language, clearly outperforming the TZK and AMP algorithms. Nevertheless, in order to avoid complexity and due to constraints imposed during its development, it can only work with words in capital letters. In the Greek language, lower case words have accent marks that can totally change the meaning of a word, like the adjectives " $\alpha \beta \alpha \theta \eta \varsigma^{\prime}$ " and " $\alpha \beta \alpha \theta \eta \zeta$ ". Both of the words mean "shallow", with the former being the masculine nominative case and the later the feminine genitive. In addition, the stemmer is able to handle suffixes but not prefixes. The essential information about all stemming algorithms that affected this thesis work and were presented in the previous two sections is listed in Table 3.

Table 3: Stemming Algorithms - Summarized Information

| Name | Langua <br> ge | Year | Web <br> Availabil <br> ity | Derivation <br> Dealing | Number of <br> Execution <br> Steps | Weaknesses |
| :---: | :---: | :---: | :---: | :---: | :---: | :--- |
| Lovins' | English | 1968 | Yes | Yes | 2 | Aggressive with <br> short stems and <br> words |
| Porter's | English | 1980 | Yes | No | 5 | Quite aggressive and <br> produces <br> overstemming <br> [Carlberger et al., <br> 2001] |
| TZK | Greek | 1995 | No | Yes | 2 | Does not handle all <br> suffixes |
| AMP | Greek | 2001 | No | No | 4 | Unable to handle <br> compound words |
| Ntais' | Greek | 2006 | Yes | No | 29 | Relatively new and <br> untested; can only <br> handle capital letters; <br> produces <br> understemming <br> errors |

### 1.6 Problems and Issues with the Latest Algorithm for Greek

According to its author [Ntais, 2006], the latest stemming algorithm for Greek suffers from a few limitations and constraints that had to be imposed during its development. One of them is the incapability to handle lower case letters. The algorithm is only able to handle words in upper case letters and will not stem any word that contains even one letter in lower case. Since words in upper case do not have tone marks in Greek, the author is solving the problem of the moving mark phenomenon that can be observed during the conjugation of verbs, nouns and adjectives since no tone mark has to be presented in the stem returned. Because of that, the algorithm has a new limitation since most of the words encountered in Greek texts are in lower case.

Another issue with the existing algorithm is incapability to process some important
suffixes. The rationale was that the inclusion of these suffixes would introduce more errors if the appropriate exception list was also introduced. The creation of an extensive exception list was not feasible at that moment, so the algorithm is not treating suffixes like "- $\alpha \tau \varepsilon$ "," $-\alpha \sigma \tau \varepsilon$ " and "- $\tau \varepsilon$ ". These suffixes correspond to words from many syntactical categories like adverbs, nouns and verbs. Verbs with these endings correspond to past tenses. Past tenses are extensively used in Greek as well as in other Mediterranean languages, since similar observations have been mentioned in the use of Spanish, Portuguese and Italian. Cultural writings and conversational contexts are widely employing tenses like past continuous. It is therefore clear, that the exclusion of these suffixes makes the algorithm incomplete and inaccurate and limits its performance.

In addition to the incompletenesses of the algorithmic design, the implementation of the algorithm offers limited usability. Ntais [2008] has provided a web interface written in Javascript. Through a form, users can insert a Greek word in upper case and have its stem returned. Despite the fact that everyone can examine the algorithm, since it is embedded in the web page, it can not be directly used by any kind of application that requires stemming in Greek. Its implementation language, Javascript, is a powerful language for client side scripting in web applications. Nevertheless, it can not be used for writing a library that can be used by other applications.

### 1.7 Aims of this Thesis

We are conducting this work in order to fully test the existing algorithm and improve it. From our initial, undocumented tests, we concluded that the existing algorithm is giving satisfactory but inaccurate and incomplete results . We are convinced that documenting and analysing the results and improving the algorithm would be a contribution not only to computer science and computational linguistics in particular but also to all these fields that Greek is used including medicine and mathematics. In addition, the algorithm can later on be used at a production level in a search engine, with the potential to give better results and a better web searching experience to users searching for documents written in Greek.

In addition to improving the existing algorithm, we will also provide a library version of the algorithm written in PHP. The reason is that PHP is currently one of the most widely used languages in the web by providing server-side scripting for building dynamic web
sites. By implementing a PHP algorithm, our aim is to provide a stemmer that can be directly used by the engine of any web application, for any kind of web search or linguistics. Other programming languages such as Javascript or even the more powerful like C and C++ lack this ability [Lerdorf and Tatroe, 2002; Flanagan, 2004]. In conclusion, our work, which will be available under an Open Source licence, will lead to a more powerful, more complete and more consistent Greek stemmer that can be directly examined, used and modified by others.

### 1.8 Research Question

In search for solutions to the previously stated problems, the research question to be answered in this thesis can be formulated as follows:

Up to which point the addition of more syntactical rules and exceptions improves the precision of the Ntais stemming algorithm?

The previous algorithm by Ntais [2006] does not, deliberately, include some suffixes in an attempt to avoid errors that occur when the appropriate exception list is not also introduced with the addition of a new suffix. The creation of an exception list is a trivial but rather time consuming process. We need to identify whether it is feasible or not to create extensive and complete lists of exceptions for new rules at this point, where the existing stemmer is already producing somewhat satisfactory results and covers most of the cases.

### 1.9 Methodology

One of the initial aims of this work was to test the original stemmer in combination with a search engine, and Google's search engine was a candidate. A web interface that would feed Google with modified, stemmed queries and unmodified ones could be easily built. The results of both modified and unmodified queries could then be compared. Unfortunately, not only the application of a stemmer in a web search engine is beyond the time limitations of this thesis, but it is also unclear whether Google is already utilizing any kind of stemming techniques for Greek. Furthermore, in a previous web search engine evaluation [Lazarinis, 2005], it is pointed out that Google returns a different number of results for different variations of the word "Athens" ( A $\theta \dot{\eta} v \alpha$ : Athens, AӨŋ́vac: of Athens, A $\theta \eta v \dot{v}$ : of (the city of) Athens). The difference in results can only imply that no stemming is used. Despite that, there are reports that some form
of stemming is being conducted [Google, 2003] although it is unclear how extensive. In addition Paice [1994] suggests that evaluating a stemmer solely in terms of IR is incomplete since IR is only one field that stemming can be applied and "...gives no insight into the specific causes of errors". Because of all these reasons and due to time limitations we decided not to test our implementations of both the existing algorithm and our improved version with a search engine.

During our research, we will modify Ntais' algorithm to use more grammatical rules, exceptions and stop words. We will improve the algorithm in a constructive and extended manner. More additions will be implemented incrementally, after testing all previous improvements each time. The task of introducing more grammatical rules is challenging since it utilises techniques and requires knowledge from two domains, computer science and linguistics. In order to evaluate Ntais' and our revised algorithm, we will execute both of them in batch mode against a collection of more than half a million Greek words. Both algorithms will stem the input words from the text, and will form groups of words that have the same stem. Our purpose is to manually check whether all the words reduced by the algorithm to the same stem also share the same semantic meaning. The two algorithms are evaluated separately and the results will be compared.

### 1.10 Overview of Thesis' Contents

In Chapter 2 we will provide a short overview of the history of the Greek language and some of the Greek grammar features and peculiarities that Greek words have during conjugation. Chapter 3 introduces some stemmer performance metrics that will be used during our evaluation in order to compare the output of the original stemmer and our modified version. The design of the existing algorithm and an extensive list of its rules are given in Chapter 4. This chapter also deals with some alternative approaches which we also considered for the re-design of our algorithm and the reasons for which we rejected them, based on their suitability to the application domain.

In Chapter 5, the evaluation of the existing algorithm is described and in Chapter 6 we describe the improvements that we decided to incorporate in our new algorithm along with the new set of rules and exception lists. The final evaluation of the algorithm and a comparison against its predecessor is given in Chapter 7.

## 2. The Greek Language

### 2.1 The History of the Greek Language

The earliest traces of written Greek can be found in more than 4400 clay tablets of the Linear B script, which was deciphered during 1951 to 1953 by the architect M. Ventris in England. This form of writing was used from 1600 to 1100 BC , and it is considered as the "earliest European script we can understand" [Robinson, 1995]. For the next 300 years, a period regarded as "the Dark Age" of illiteracy in Greece [Robinson, 1995], no traits of the Greek language have been discovered. During this period, the Homeric Greeks gave their position to the classical Greeks. The classic period of Ancient Greece (500-323 BC) coincides with the emergence of a new alphabet borrowed from the Phoenicians. Although it is debatable whether the Phoenicians or Greeks living in Phoenicia were the creators of this alphabet [Robinson, 1995], this alphabet is the ancestor of not only the Greek alphabet, but through the Etruscan and Latin languages, the ancestor of modern European alphabets.[Baugh \& Cable, 2001]. The known fact is that the first consonant-only based Phoenician alphabet came to Greece without vowels, and the ancient Greeks added vowels to it. These added letter-characters improved greatly the communication and increased the use of the alphabet in everyday life and in writing form. After all, the enhanced with vowels new alphabet came nearer to the needs of everyday speech and it mirrored the spoken words more clearly and more accurately than its consonant based predecessor.

After many additions and simplifications through the years, the nowadays alphabet contains twenty four (24) letters. Table 4 presents the latest form of the Greek alphabet including both upper and lower case letters and their equivalent sound in the English alphabet. The third character of Sigma ( $\varsigma$ ) is only used in the end of a lower case word.

Table 4: The Greek Alphabet. Letters and their equivalent sound in English

| Alpha | Beta | Gamma | Delta | Epsilo <br> n | Zeta | Eta | Theta | Iota | Kapa | Lamda | Mi |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| A $\alpha$ | B $\beta$ | $\Gamma \gamma$ | $\Delta \delta$ | $\mathrm{E} \varepsilon$ | $\mathrm{Z} \zeta$ | $\mathrm{H} \eta$ | $\Theta \theta$ | I 1 | K к | $\Lambda \lambda$ | $\mathrm{M} \mu$ |
| a | V | - | $\begin{gathered} \text { th } \\ \text { (the) } \end{gathered}$ | e | Z | 1 | th | i | k | 1 | m |
| Ni | Xi | Omikron | Pi | Ro | Sigma | Tau | Ypsilon | Phi | Chi | Psi | Omega |
| Nv | $\Xi \xi$ | O | $\Pi \pi$ | P $\rho$ | $\Sigma \sigma \varsigma$ | $\mathrm{T} \tau$ | Y $v$ | $\Phi \varphi$ | X $\chi$ | $\Psi \psi$ | $\Omega \omega$ |
| n | ks | 0 | p | r | S | t | i | f | h | ps | 0 |

The Ancient Greeks of the Classical period were organized in city states. Each of the main city states had its own region of influence and with that a different dialect. The differences between these dialects were minor, so Greek was considered as a common language [Triantafyllidis, 1941]. It was only after the conquests of Alexander the Great in Asia when the Athenian dialect, after borrowing words from other dialects, became the common dialect spoken from Greece and Egypt to Syria and Persia. The language continued to evolve for the next centuries, until the fall of Constantinople in 1453 and the beginning of the Ottoman era. For the next 400 years, following the closing of schools, the Greek language is kept oral, divided into local dialects. A few examples of written material from this period can be credited to Greeks living in countries in Central and Western Europe, like Romania, Austria, Russia, the Hungarian Empire and Italy.

After the Greek War of Independence, which started in 1821, and the liberation in 1829, two competing varieties are found. The popular and spoken "Dimotiki", used among the people, and the official and most resembling to Ancient Greek "Katharevousa", used mostly by the intellectuals. Today, modern mainstream Greek is based on "Dimotiki" and is the official state language, since 1975, with simplifications in the intonation
system since 1981. In places, there are still local dialects with varying degrees of differentiation from the mainstream language.

Greek is spoken by 14 to 17 million people, officially in Greece and Cyprus and unofficially in countries like Australia, USA and Canada, where there are Greek and Cypriot communities. In addition, medical and philosophical terms are often Greek, Greek-derived or combinations of Latin and Greek words [Kurz \& Kilian, 2001]. The domain of Humanities is undeniably a language world with terms and concepts that have originally been founded in the subject of Philosophy and Mathematics and expressed in Greek.

### 2.2 Stemming in Greek

The Greek language is grammatically more complex than English. It has conjugations and morphologically complex words [Mackridge, 1987]. Articles, adjectives, nouns and even first names and surnames may be in various cases (like nominative and genitive), in singular or plural form and they are differentiated according to their gender (masculine, feminine, neuter). Table 5 contains the singular and plural numbers of all cases and genders that the word "cat" that be found in Greek.

Table 5: The noun "cat" in Greek

| Singular |  |  |  | Plural |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Cases | masculine | feminine | neuter | masculine | feminine | neuter |
| nominative | $\gamma \alpha ́ \tau 0 ¢$ | $\gamma$ ¢о́ $\boldsymbol{\alpha}$ | $\gamma \alpha \tau i$ | $\gamma \alpha ́ \tau 0!$ | $\gamma \alpha \dot{\tau}$ ¢¢¢ | $\gamma \alpha \tau \iota \alpha ́$ |
| genitive | $\gamma \alpha ́ \tau 0 v$ | $\gamma \alpha{ }^{\prime} \tau \alpha$ | $\gamma \alpha \tau$ ¢0v | $\gamma \alpha ́ \tau \omega v$ | $\gamma \alpha ́ \tau \omega v$ | $\gamma \alpha \tau \iota \omega ்$ |
| accusative | $\gamma \alpha{ }^{\text {co }}$ | $\gamma \boldsymbol{\chi} \boldsymbol{\alpha} \boldsymbol{\alpha}$ | $\gamma \alpha \tau i$ | үótovs | $\gamma \alpha \dot{\tau}$ ¢ | $\gamma \alpha \tau \iota \alpha{ }^{\prime}$ |
| vocative | $\gamma \alpha ́ \tau \varepsilon$ | $\gamma$ ¢о́ $\boldsymbol{\alpha}$ | $\gamma \alpha \tau i$ | $\gamma \alpha ́ \tau 0!$ | $\gamma \alpha \dot{\tau}$ ¢¢¢ | $\gamma \alpha \tau \iota \alpha ́$ |

In addition to the complexity mentioned above, verbs are also heavily inflected. The

Greek language consists of present, past and future tenses with both perfective and imperfective aspects. These tenses have active and passive voices. Verbs are divided into two conjugations classes which have different endings and many times there are alternative endings for the same number and person. From the tables given in Appendix A, containing some examples about verb conjugation, it is obvious that Greek is much more complicated than English. Where in English four (4) endings are used, in Greek the distinct endings are 107, even without counting the different endings because of the moving mark phenomenon. Not only a stemmer has to deal with an enormously greater number of endings, but a somewhat perfect stemmer should be aware of how to deal with the "moving" tone mark issue.

## 3. Stemmer Performance Metrics

In order to evaluate the existing stemmer and measure its effectiveness, we will introduce some of the metrics that can be found in the previous literature.

### 3.1 Frakes' Metrics

Frakes [2003] defines stemmer strength as the degree to which a stemmer changes words. These changes fall into two categories, removal and recording. Removal is the decrease of a word's length due to elimination of an affix, whereas recording is the replacing of a word's letter with another. Since the strength of a stemmer can affect the precision and recall in queries, Frakes defines a set of metrics that help to compare algorithms by having the algorithms stem the same texts and compare the results of the following metrics:

- The Mean Number of Words per Conflation Class

The mean number of words per conflation class is the average number of words that are found in each conflation class. For example if the words "engineer," "engineered," and "engineering" are stemmed to "engineer," then this conflation class size is three. Stronger stemmers produce conflation classes with more words than lighter stemmers, from the same text.

- Index Compression Factor

The index compression factor is defined as ICS $=\frac{n-s}{n}$, where $n$ is the number of words in the corpus and $s$ is the number of stems produced by the stemmer. This metric indicates the index reduction that can be achieved through stemming. For example, if during the stemming of a corpus with 1000 words ( $n$ ) we end up with 800 stems ( $s$ ), we have eliminated 200 words which means an index compression factor of $20 \%$. Stronger stemmers will have a larger index compression factor than lighter stemmers.

- The Number of Words and Stems that Differ

Stemmers often leave words unchanged. The reason behind this behaviour can be either the lack of an appropriate rule, a software bug in the implementation of the algorithm or a design choice from the authors. For example, a stemmer might not alter "engineer" because it is already a dictionary entry. A big ratio of unchanged words to total words can indicate poor algorithmic performance. Stronger stemmers will change words more
often than weaker stemmers.

- The median and mean modified Hamming distance

The Hamming distance between two strings of equal length is defined as the number of characters in the two strings that are different at the same position. For strings of unequal length we add the difference in length to the Hamming distance to give a modified Hamming distance function d. This measure takes into account transformations of stem endings. For example, a stemming algorithm might reduce the corpus $\{$ try, tried, trying \} to the stem "tri". The mean modified Hamming distance between the original words and the stem is $\mathrm{D}=(1+2+4) / 3=2.33$ characters, and the median is 2 .

### 3.2 Error Metrics

There are two clearly distinct error metrics categories concerning stemmers, understemming and overstemming.

Understemming occurs when words are not fully stemmed to their potential stem. In that case, words that share the same conceptual meaning are stemmed to different stems and assigned to a different conflation class.

In contrast, overstemming occurs when words that do not share the same conceptual meaning are reduced into the same stem and assigned to the same conflation class. According to other evaluations [Alvares et al., 2005], the most accurate way to check for understemming and overstemming errors is through human interference. Some examples for both categories are given in Table 6. The first example shows understemming where two words have different stems although they should had the same, since they both have to do with "selling". On the other hand, the second examples demonstrates overstemming where two words with different semantics, "selling " and "bird", are incorrectly reduced to the same stem.

## Table 6: Stemming Errors

| Word | Meaning | Stem |
| :---: | :---: | :---: |
|  | （I）sell | $\pi$ ои入－ |
| $\pi$ тол入ө́vtas | selling | $\pi$ тои入ovт－ |
| Overstemming |  |  |
| Word | Meaning | Stem |
| $\pi$ тол $\alpha^{\prime} \omega$ | I sell | тоטג－ |
| тоט入í | bird | лоט入－ |

## 4. Algorithmic Design

### 4.1 The Existing Algorithm

The algorithm provided by Ntais deals "with each suffix individually" in a decentralized manner [Ntais, 2006]. The algorithm has 29 rules that treat 158 suffixes. Every rule is executed in an individual step and a set of suffixes is provided in order to remove the longest matching suffix. In all but the first steps, a list of exceptions is also examined and some different suffixes are added to the stem if needed in order to deal with the complexity of the Greek language. Additionally, each step may have its own exceptions. We have decided to keep this design and base our work on this.

Table 7 presents the algorithm by Ntais [2008]. In each step the rule with suffixes to be examined, along with actions to be taken and exceptions to be considered are described.

Table 7: Ntais' Algorithm

| Step \# | Rule | Action | Exceptions |
| :---: | :---: | :---: | :---: |
| 1 | Word ends in: <br> ФАГІА\|ФАГІОҮ| <br> ФАГIQN\|इКАГIA| <br> ГКАГIOY\|ЕКАГIQN| <br> ОАОГІОҮ\|ОАОГІА| <br> OAOГISN\|EOFIOY| <br> इOГIA\|इOГIQN| <br> TATOГIA\|TATOГIOY| <br> TATOFIQN\|KPEAइ| <br> KPEATOL\|KPEATA| <br> KPEATQN\| ${ }^{\text {KEPAE\| }}$ <br> ПEPATOL\|ПEPATA| <br> IEPATSN\|TEPAE| <br> TEPATOL\|TEPATA| <br> TEPATQN\|T $\Omega \Sigma\|\Phi \Omega T O \Sigma\|$ <br>  <br> KAQELTRE\| <br> KAQEETRTOL\| <br> KAGELTRTA\| <br> KA@ELTRTSN\| <br> ГEГONOE\| <br> ГETONOTOE\| <br> ГEГONOTA\| <br> ГEГONOTAN | Replace suffix with: <br> $\boldsymbol{\Phi A}\|\Sigma K A\|$ <br> OAO\|EO| <br> TATO\| <br> KPE\|ПEP| <br> $T E P\|\Phi \Omega\|$ <br> KA $\mathcal{E E \Sigma T \|}$ <br> ГЕГОN |  |


| 2a | Word ends in： AUEL｜AASN | Remove suffix／ Check exceptions | If after removal the word does not end in： <br> OK｜MAM｜MAN｜MПAMI｜ <br> ПATEP $\|\Gamma I A \Gamma\|\|N T A N T\| K Y P\|\Theta E I\|$ ПЕЄЕР <br> Add＂ $\boldsymbol{A} \boldsymbol{u}^{\prime}$＂in the end |
| :---: | :---: | :---: | :---: |
| 2b | Word ends in： ELEL｜ELSN | Remove suffix／ Check exceptions | If after removal the word ends in： <br> ОП｜ІП｜ЕМП｜ҮП｜ГНП｜ <br> $\triangle$ АП｜КРАГП｜МІ $\Lambda$ <br> Add＂EA＂in the end |
| 2c | Word ends in： OYAES｜OYASN | Remove suffix／ Check exceptions | If after removal the word ends in： АРК｜КАИІАК｜ПЕТАИ｜ИIX，｜ ПИЕХ $\|\Sigma K\| \Sigma\|\Phi \Lambda\| \Phi P \mid$ ВЕ $\mid$｜ИOY $\mid$ $X N\|\Sigma \Pi\| T P A \mid \Phi E$ <br> Add＂OY4＂in the end |
| 2d | Word ends in： E $\Omega \Sigma \mid E \Omega N$ | Remove suffix／ Check exceptions | If after removal the word is one of： АРК｜КАЛІАК｜ПЕТАА｜ИIX｜ ПИЕЕ $X N\|\Sigma \Pi\| T P A \Gamma \mid \Phi E$ <br> Add＂$E$＂in the end |
| 3 | Word ends in： IA｜IOY｜ISN | Remove suffix／ Check exceptions | If after removal the word ends in Vowel <br> Add＂$\Gamma$＂in the end |
| 4 | Word ends in： IKA｜IKO｜IKOY｜IKSN | Remove suffix／ Check exceptions | If after removal the word ends in Vowel <br> OR <br> is one of ： <br> A $A\|A \Delta\| E N A\|A M A N\|$ <br> AMMOXA $\|\mathbf{H \Theta}\| A N H \Theta \mid$ <br> ANTIU｜ФYУ｜BPSM｜ГEP｜ <br> EEST｜KA1П｜KAMAIN｜ <br> KATAL｜MOYイ｜MПAN｜ <br> МПАГІАТ｜МПОА｜МПОГ｜ <br> NIT｜EIK｜工YNOMHИ｜ <br> ПЕТЕ｜ПITE｜IIKANT $\mid$ <br> ПАIATE｜ПOГTEAN｜ <br> IIPSTOU｜EEPT｜EYNAU｜ <br> TEAM｜YПOU｜TINON｜ <br> ФYイOU｜XAE <br> Add＂ $\boldsymbol{I}$＂in the end |


| 5a | Word is $\boldsymbol{A C A M E}$ | Stem is AГАМ |  |
| :---: | :---: | :---: | :---: |
| 5a | Word ends in : <br> AГAME $\mid$ HEAME <br> OYEAME\|H*KAME $\mid$ <br> H@HKAME | Remove suffix / Check exceptions |  |
| 5a | Word ends in: $A M E$ | Remove suffix / Check exceptions | If after removal the word is one of: ANAП\|AПOQ|AПOK|AПOET| ВОYВ|ЕЕЄ|ОУ^|ПЕЄ|ПIKP| ПOT $\|\Sigma I X\| X$ <br> Add " $\boldsymbol{A} \boldsymbol{M}$ " in the end |
| 5b | Word ends in: <br> ACANE\|HEANE| OYEANE|IONTANE| IOTANE|IOYNTANE| ONTANE $\mid$ OTANE $\mid$ OYNTANE\|HKANE| HOHKANE | Remove suffix / Check exceptions | If after removal the word one of: $T P \mid T \Sigma$ <br> Add "AГAN" |


| 5b | Word ends in: <br> AFANE\|HEANE <br> OYEANE\|IONTANE| <br> IOTANE\|IOYNTANE| <br> ONTANE $\|O T A N E\|$ <br> OYNTANE\|HKANE| <br> H@HKANE | Remove suffix / Check exceptions | $\square$ If after removal the word ends in Vowel without " $\boldsymbol{Y}$ " <br> OR <br> is one of : <br> BETEP $\mid$ BOYAK $\mid$ BPAXM $\mid \Gamma\rceil$ ДРААОYM\|O|КААПОYZ| <br> КАЕТЕИ\|КОРМОР|ИАОПИ| <br> M@AME@\|M|MOYEOYAM|N| <br> ОҮИ\|П|ПЕИЕК|ПИ|ПОЛІІ| <br> ПОРТОА\|ЕAPAKATE|EOYAT] <br> TEAP1AT\|OPФ|TEIГГ|TEOП| <br> ФЛТОЕТЕФ\|Х|ЧҮХОПИ|AГ| <br> ОРФ\|ГАИ|ГЕР|ЗЕК|ЗІПИ| <br> AMEPIKAN $\mid$ OYP $\mid$ IIU@\|IOYPIT $\mid$ <br> $\Sigma\|Z \Omega N T\| I K\|K A \Sigma T\| K O \Pi\|\Lambda I X\|$ <br> תOY@HP\|MAINT|ME1| $\Sigma I \Gamma\|\Sigma \Pi\|$ <br> $\Sigma T E \Gamma\|T P A \Gamma\| T \Sigma A \Gamma\|\Phi\| E P\|A \Delta A \Pi\|$ <br> A@IГГ $\|A M H X\| A N I K\|A N O P \Gamma\|$ <br> АПНГ $\mid$ АПIӨ\|АТГІГГ $\mid$ ВА $\Sigma \mid$ <br> ВАГK $\mid$ BA@YГA $\mid$ BIOMHX $\mid$ <br> BPAXYK $\mid$ IIIAT $\mid$ IIAT\|ENOP $\mid$ <br> $\Theta Y \Sigma \mid$ KAПNOBIOMHX <br> КАТАГАИ $\mid$ КАІВ $\mid$ КОІІАРФ\|ИIB $\mid$ <br> МЕГІОВІОМНХ <br> MIKPOBIOMHX $\mid$ NTAB $\mid$ <br> ЕНРОКАІВ\|ОАІГОААМ| <br> ОАОГАА\|ПЕКТАРФ|ПЕРНФ| <br> ПЕРІТР\|ПААТ|ПОЛУААП| <br> ПОАУМНХ $\|\Sigma T E \Phi\|$ TAB $\mid$ TET $\mid$ <br> ҮПЕРНФ\|УПОКОП| <br> ХАМНАОДАП\|ЧНИОТАВ <br> Add " $\boldsymbol{A} \boldsymbol{N}$ " in the end |
| :---: | :---: | :---: | :---: |
| 5c | Word ends in: HEETE | Remove suffix / Check exceptions |  |


| 5c | Word ends in: ETE | Remove suffix / Check exceptions | If after removal the word ends in Vowel without " $P$ " <br> OR <br> is one of : <br> ABAP $\|B E N\| E N A P\|A B P\| A \Delta\|A \Theta\|$ AN $\mid$ AПИ\|BAPON $\mid$ NTP\| $\Sigma К \mid$ KOП $\mid$ МПОР\|NIФ|ПАГ|ПАРАКАИ| $\Sigma E P \Pi\|\Sigma K E \Lambda\| \Sigma Y P \Phi\|T O K\| Y\|\Lambda\| E M \mid$ $\boldsymbol{\Theta A P P \|} \mid \boldsymbol{\Theta}$ <br> OR <br> ends in : <br>  EYP $\mid$ TIQ $\mid$ VПEP@ $\mid$ PA $\Theta\|E N \Theta\|$ $P O \Theta\|\Sigma \Theta\| \Pi Y P \mid A I N] \Sigma Y N \Delta \mid \Sigma Y N]$ $\Sigma Y N \Theta\|X \Omega P\| \Pi O N\|B P\| K A \Theta\|E Y \Theta\|$ EKQ\|NET|PON|APK|BAP|BOA| תФЕ $\boldsymbol{\Omega}$ <br> Add " $\boldsymbol{E T}$ " in the end |
| :---: | :---: | :---: | :---: |
| 5d | Word ends in: ONTAE\| $\boldsymbol{\Omega N T A \Sigma}$ | Remove suffix / Check exceptions | If after removal the word is: APX <br> add "ONT" in the end OR <br> If after removal the word is: KPE <br> add " $\boldsymbol{\Omega} \boldsymbol{N} \boldsymbol{T}$ " in the end |
| 5e | Word ends in: OMAETE\|IOMAETE | Remove suffix / Check exceptions | If after removal the word is: ON <br> add "OMAET" in the end |
| 5 f | Word ends in: IELTE | Remove suffix / Check exceptions | If after removal the word is one of: $\boldsymbol{\Pi}\|A \Pi\| \Sigma Y M \Pi\|A \Sigma Y M \Pi\| A K A T A \Pi \mid$ AMETAMФ <br> Add "IELT" in the end |
| 5 f | Word ends in: EETE | Remove suffix / Check exceptions | If after removal the word is one of: <br>  AP\|IPO|NIL <br> Add "ELT" in the end |
| 5g | Word ends in: H@HKA\|H@HKE H@HKE | Remove suffix / Check exceptions |  |


| 5 g | Word ends in: HKA\|HKE $\mid$ \|HKE | Remove suffix / Check exceptions | If after removal the word is one of: <br> $\triangle I A \Theta\|\Theta\| \Pi A P A K A T A \Theta\|\Pi P O \Sigma \Theta\|$ $\Sigma Y N \Theta$ <br> OR <br> ends in: <br> $\Sigma K \Omega \Lambda\|\Sigma K O Y \Lambda\| N A P \Theta\|\Sigma \Phi\| O \Theta \mid$ <br> $\boldsymbol{\Pi I \Theta}$ Add "HK" in the end |
| :---: | :---: | :---: | :---: |
| 5h | Word ends in: <br> OYEA\|OYEE $\mid$ OY $E$ | Remove suffix / Check exceptions | $\square$ If after removal the word is one of: <br> ФAPMAK\|XAA|AГK|ANAPP| <br>  ПAT $\|P\| \Lambda\|M E \Delta\| M E \Sigma A Z \mid$ <br> YIIOTEIN\|AM|AIO|ANHK| <br> $\triangle E \Sigma I O Z\|E N \triangle I A \Phi E P\| \triangle E \mid$ <br> $\triangle E Y T E P E Y\|K A \Theta A P E Y\| \Pi \Lambda E \mid T \Sigma A$ <br> OR <br> ends in: <br> ПОААР\|ВАЕП|ПANTAX|ФРУム| MANTI 1 |MAム 1 |KYMAT $\mid$ AAX $\mid$ АНГ $\|\Phi А Г\| О М \mid П Р \Omega Т ~$ Add "OYI" in the end |


| 5 i | Word ends in: <br> $\mathbf{A Г A}\|\mathbf{A Г E} \mathbf{\Sigma}\| \mathbf{A Г E}$ | Remove suffix / Check exceptions | If after removal the word is one of: <br> АВАГТ\|ПОАҮФ|АВНФ|ПАМФ| <br>  ANYET $\mid$ AПEP $\mid$ AГПAP\|AXAP| पEPBEN| $\triangle$ POГOП\|ЕEФ|NEOI | NOMOT|OAOI|OMOT|IPOET| ПРОГЛПОП|ธYMП| $\Sigma Y N T\|T\|$ YПOT\|XAP|AEIIT|AIMOET $\mid$ ANYI\|AПOT|APTII|UIAT|EN| ЕПIT|КРОКААОП|ЕІАНРОП|ム| NAY|OY1AM|OYP|ITTP|M <br> OR <br> ends in: <br>  ПP\|LOX $\mid \Sigma M H N$ <br> BUT <br> is not one of: <br> $\boldsymbol{\Psi O T} \mid$ NAYIOX <br> AND <br> does not end in: <br> KOM <br> Add " $\boldsymbol{A} \boldsymbol{\Gamma}$ " in the end |
| :---: | :---: | :---: | :---: |
| 5 j | Word ends in: $H \Sigma E\|H \Sigma O Y\| H \Sigma A$ | Remove suffix | If after removal the word is one of: N\|XEPEON| $4 \Omega \Delta E K A N \mid$ EPHMON\|MEГAMON|EПTAN <br> Add " $\boldsymbol{H} \boldsymbol{\Sigma}$ " in the end |
| 5k | Word ends in: HETE | Remove suffix / Check exceptions | If after removal the word is one of: $A \Sigma B\|\Sigma B\| A X P\|X P\| A \Pi \Lambda\|A E I M N\|$ पY $\mathbf{I X P} \mid$ EYXP\|KOINOXP $\mid$ ПААIMY <br> Add " $\boldsymbol{H \Sigma T}$ " in the end |
| 51 | Word ends in: <br> OYNE\|HEOYNE| <br> H@OYNE | Remove suffix / Check exceptions $\square$ | If after removal the word is one of: N\|P| $\Sigma I I \mid$ ITPABOMOYT $\mid$ KAKOMOYTE\|EESN <br> Add " $\boldsymbol{O Y} \boldsymbol{N}$ " in the end |


| 51 | Word ends in: OYME\|HEOYME H@OYME | Remove suffix / Check exceptions $\square$ | If after removal the word is one of: ПАРАГОYГ $\|\Phi\| X\|\Omega P I O \Pi \Lambda\| A Z \mid$ AMAOEOYE\|AEOYE <br> Add "OYM" in the end |  |
| :---: | :---: | :---: | :---: | :---: |
| 6 | Word ends in: <br> MATA\|MATSN|MATOE | Remove suffix / Check exceptions | Always <br> Add " $\boldsymbol{M} \boldsymbol{A}$ " in the end |  |
| 6 | Word ends in: <br> A\|ACATE $\mid$ AIFAN $\mid$ AEI $\mid$ AMAI $\mid$ AN $\mid$ A $\Sigma\|A \Sigma A I\| A T A I\|A \Omega\| E\|E I\| E I \Sigma \mid$ <br>  <br>  IONTOY氵AN\|IOEAETAN|IOEAETE|IOEOYN|IOEOYNA| IOTAN|IOYMA|IOYMA ETE|IOYNTAI|IOYNTAN|H|HUE $\mid$ HUSN\|H@EI|H@EIL|H@EITE|H@HKATE|H@HKAN| <br>  HEE $\Sigma\|H \Sigma O Y N\| H \Sigma \Omega\|O\| O I\|O M A I\| O M A \Sigma T A N\|O M O Y N\|$ OMOYNA\|ONTAI|ONTAN|ONTOYइAN|Oइ|OEAETAN| OEAETE|OEOYN|OEOYNA|OTAN|OY|OYMAI|OYMAETE| OYN $\mid$ OYNTAI\|OYNTAN|OY $\mid$ \|OYEAN $\|O Y \Sigma A T E\| Y\|Y \Sigma\| \Omega \mid \Omega N$ |  |  | Remove suffix |
| 7 | Word ends in: <br> ELTEP\|ELTAT $\mid$ OTEP $\mid$ <br> OTAT\|YTEP|YTAT|STEP| תTAT | Remove suffix |  |  |

### 4.2 Rejected Designs

During the evaluation of the existing work we rejected some alternative to Ntais' algorithm designs for our stemmer.

### 4.2.1 Lovins' Design

As mentioned in Section 1.4, Lovins' algorithm uses two steps in order to remove suffixes from words. The algorithm is used for stemming in English. As a design, it offers more simplicity than the algorithm provided by Ntais, which is executed in 29 steps. During our evaluation, we implemented a design similar to Lovins' algorithm using the same list of suffixes that Ntais used. The reason was that we observed Ntais algorithm's behaviour and it removes suffixes in one or two steps on average, despite the fact that it always executes on 29 steps. Unfortunately, our implementation of a Lovinslike algorithm didn't offer any improvement. Ntais' algorithm uses an exception list after each step in order to deal with the richness and irregularities of Greek. Even if 29 steps may be regarded as poor design, since the rest of the algorithms mentioned in Table 3
execute in two to five steps, tracking and matching endings with exceptions is easily conducted. Thus, the algorithm can be studied and improved easily. Despite of the fact that eventually in our algorithm we kept Ntais' design and we even added more steps, we made sure that the algorithm is not making unnecessary checks by returning the correct stem right after matching all possible suffixes.

### 4.2.2 Context-Free Grammar Design

Another alternative design that we examined was a more theoretical and sophisticated one based on the theory of context-free grammars. The main idea behind the theory is that all natural languages are based on elements, which in turn can be based into other elements recursively [Lewis and Papadimitriou, 1998]. Parse trees can be used and by applying rewrite grammatical rules the elements of sentences can be constructed and, in the case of a stemmer, de-constructed and analysed.

This approach is similar to corpus-based stemming, where each word is examined and assigned to a conflation class not only according to the grammatical meaning it holds but also to its semantic meaning it contains in the text [Xu and Croft, 2000]. Our hypothesis was that an algorithm based in these principles could be trained to recognize words and build more complex conflation classes. In that way, when a word was given, the stemmer could return a list of alternative words that have already been extracted from texts and have the same meaning. Nevertheless, this approach is beyond stemming even though it could, theoretically, offer better results in search queries. Another prohibitive reason was that although there is literature available about this subject, there is no evidence, to our knowledge, of any production or working systems and stemmer implementation using this approach, at least publicly available.

### 4.2.3 Dictionary Based Design

A rule based stemming algorithm has to be aware of exceptions. By adding more rules to an algorithm, a researcher must also add more exceptions, in order to deal with the complexity of a natural language. One may argue that a stemmer with many exceptions resembles a dictionary stemmer, since it has embedded in its design lists of words. A dictionary stemmer could be built by using a rule-based stemmer for creating an initial list of stems, and then manually checking the stems created. We rejected that design, since our personal interest was not only to create a solution for Greek stemming but also to study the behaviour of the language.

### 4.2.4 Krovetz's Experimental Algorithm

Another possible use of a dictionary for our algorithmic design could be using a dictionary to check whether a rule produced a word that is a dictionary entry. Krovetz [1993] experimented with this design. In his algorithm, prior to the execution of each rule, the word is looked up in a dictionary. If the word is an entry in the dictionary, the algorithm returns it. In the opposite case, the algorithm proceeds in evaluating more rules. When the word is modified by a rule, the resulting word is again looked up in the dictionary. This pattern continues until the remaining word is an entry or no more rules can be applied.

Krovetz experimented with this design, in an attempt to deal with the aggressiveness of Porter's algorithm and ended up with a "weaker "stemmer. For example, this experimental algorithm would stem "generalisations" to "generalisation" and not to "general". This would provide optimal results in IR since "general" is a word with multiple entries in a dictionary. One of these entries has the same semantic meaning with "generalisations" while the other deals with military ranks. A search engine using Porter's algorithm would probably retrieve documents of both categories while trying to serve a query with the term "generalisation".

We decided not to follow this design for three reasons. Our main objection in following this design is the fact that we are conducting this work having both linguistics and IR in mind. A stemming algorithm with similar design could possibly produce better results in IR, but it would be an incomplete algorithm. In addition, this design would require many modifications in the suffix and exception lists. Krovetz's design is manipulating suffixes in a piece by piece fashion. In contrast, Ntais' algorithm is removing suffixes in one or two steps. By following Krovetz's approach, the suffix and exception list would have to be created from the beginning, a task that would require an enormous amount of effort. Finally, Krovetz reported [1993] that in many cases this design offered poor results compared to the original algorithm by Porter, even in IR uses.

### 4.3 Evaluation of Ntais' Algorithm

In order to evaluate the algorithm developed by George Ntais, we first ported the algorithm in PHP. Furthermore, we implemented a set of helper applications that will be using directly the algorithm and will keep statistics about the stems returned. The operating system used for the evaluation was Gentoo Linux but because of the
portability of PHP and our style of coding the source code is portable and can be used in any platform that PHP is ported to.

Our evaluation begun by executing our port of Ntais' stemmer against a list of Greek words. We set the application in a manner that words with co mmon stems will be grouped into conflation classes and then the output will be directed into a text file. This text file was examined manually for understemming and overstemming errors. In addition to that, we used modified Hamming Distances in order to find similar stems. From our observations, we concluded that two stems with a modified Hamming Distance of four or less can be possibly merged into one, indicating an understemming or overstemming error of the stemmer that generated them. An example is the comparison of the "BA $\Delta \mathrm{IZ}$ " and "BA $\Delta \mathrm{IZAT}$ "stems, erroneously generated by the original stemmer instead of one. The first is a stem for 23 words that have to do with the verb "to walk" while the latter is the stem generated for the word "BA $\triangle$ IZATE" which is the third person plural in the Past Continuous tense of the same verb. The modified Hamming Distance between them is 2 . We implemented a small application that outputs pairs of stems generated by the stemmer with a modified Hamming Distance between them of 3 or less. The pairs generated were potential understemming or overstemming errors and candidates for merge into one conflation class. After thorough examination of the lists generated we begun with the modification and the improvement of the algorithm.

## 5. Our Improvements

### 5.1 Introduction of Stop-Word Elimination

Stop-word removal is one of the most commonly used techniques in IR [Baeza-Yates and Ribeiro-Neto, 1999; Risvik et al., 2003; Lazarinis, 2007]. We use stop-word elimination in order to improve the performance of the stemming algorithm. The stopword list mainly contains words of length of at most four letters. These words are mainly articles, adverbs and conjunctions that can not be conjugated or stemmed. In addition, some common words that can be found in Greek texts, like initials, are also added to this list. In our modified algorithm, stop-word elimination is the first step of execution, a step that not only produces better results but also improves the running time of the algorithm. These words tend to deceive the stemmer and as a result, non existing stems of minimal length are created. The initial algorithm by Ntais solved this problem by processing only words of four letters or more. Although this approach left just a few words of 3 that could be stemmed unprocessed, we decided to add a stop-word list of more than 500 words, in order to increase precision.

### 5.2 Addition of More Grammatical Rules

In the course of our evaluation, and in order to conduct our algorithm testing and comparison, we created a set of helper applications that directly use our implementations of both Ntais' and our modified algorithm. One of these applications uses as input a list of words and creates conflation classes according to the stems returned by the stemmers. These classes were manually checked for overstemming and understemming errors in a manner similar to previous literature [Alvares et al., 2005; Ntais, 2006]. According to the results, more suffixes were added in order to deal with understemming. As Ntais [2006] had pointed out, the introduction of more rules for additional suffixes raises stemming errors due to overstemming. In order to deal with this situation, we additionally added more exceptions in order to deal with overstemming and keep precision at an acceptable level. Our efforts concentrated mainly on the the addition of rules that deal with past tenses as they play a great role and can be often found in Greek texts.

### 5.3 Introduction of Lower Case Letters

The initial stemming algorithm of Ntais only accepts as input words in upper case letters, as we mentioned in Section 1.6. Our algorithm is capable of handling words given in any case, upper, lower or combinations of both. The main body of the algorithm remains unchanged and all rules are still in capital letters. Before the evaluation of the rules, each letter in the given word is converted into upper case. The algorithm stores the case of each letter individually in a different variable. The algorithm examines all rules and creates the stem in upper case. Before returning the stem of the given word, a final alteration of the stem occurs as the algorithm consults the case of each letter on the stem, and alters the case of a letter if needed.

The initial algorithm was only accepting words in upper case in order to deal with the "moving" tone mark. Although in our implementation the problem of the "moving" tone-mark still remains, we decided to treat both upper case and lower case words. The complexity of the Greek language and the time limitations of this thesis prohibit any serious attempt to solve this problem. Moreover, in a hypothetical situation in which this problem was solved, the improvement in precision would be minimal. Finally we believe that since a stem is not a real word, but just a linguistic unit, returning stems in Greek without tone marks in not an important issue.

### 5.4 The Revised Algorithm

After careful examination of the output of the original stemmer, we tried to incorporate as many modifications as possible. The original algorithm has some inaccuracies but, searching for omissions and errors can be compared with searching for a needle in a haystack. Moreover, an addition of a rule that corrects some errors may create other errors, unless an appropriate exception list is also created. Nevertheless, we added more rules in order to correct wrong patterns that kept appearing in the output. One striking example was the omission of any rules for suffixes that appear in Past Continuous (IZA, IZEE, IZE, IZAME, IZATE IZAN) and past tenses in general. This detailed work appears in Table 8.

Table 8: Our revised algorithm (additions and modifications highlighted)

| Step \# | Rule | Action | Exceptions |
| :--- | :--- | :--- | :--- |
| UL1 | $\begin{array}{l}\text { Word contains letters in } \\ \text { lower case }\end{array}$ | $\begin{array}{l}\text { Convert letters in upper case } \\ \text { Store their position in the word }\end{array}$ |  |
| SWR | $\begin{array}{l}\text { Word is one of the Stop } \\ \text { Word List (Apendix C) }\end{array}$ | $\begin{array}{l}\text { Return word } \\ \text { unchanged }\end{array}$ |  |
| 1 | Word ends in: | Replace |  |
| suffix with: |  |  |  |$]$


| S1 | Word ends in: <br> IZA\|IZE $\mid$ IZE $\mid$ IZAME $\mid$ <br> IZATE\|IZAN|IZANE| <br> IZQ\|IZEIE|IZEI| <br> IZOYME\|IZETE|IZOYN| <br> IZOYNE | Remove suffix / Check exceptions / Exit | If after removal the word ends in: <br> ANAMПA\|EMПA|EПA <br> ЕАNАПА\|ПА|ПЕРІІА $\mid$ A $\Theta P O \mid$ <br> $\Sigma Y N A \Theta P O \mid \triangle A N E$ <br> Add " $l$ " in the end <br> If after removal the word ends in: <br> MAPK\|KOPN|AMIIAP|APP| <br> ВАЄYPI\|BAPK|B|BOABOP|ГКР| <br> ГАУКОР\|ГАУКУР|ІМП|И|ИОҮ| <br> МАР\|М|ПР|МПР|ПОАУР|П|Р| <br> ПІПЕРОР <br> Add " $I Z$ " in the end |
| :---: | :---: | :---: | :---: |
| S2 | Word ends in: <br> $\Omega \Theta H K A\|\Omega \Theta H K E \Sigma\|$ <br> 』ӨHKE $\|\Omega \Theta H K A M E\|$ <br> ЛఆHKATE $\|\Omega \Theta H K A N\|$ <br> 』@HKANE | Remove suffix / Check exceptions / Exit | If after removal the word is: $\boldsymbol{A} \boldsymbol{\Lambda}\|\boldsymbol{B I}\| \boldsymbol{E N}\|\boldsymbol{Y \Psi}\| \mathbf{I I}\|\boldsymbol{Z} \boldsymbol{\Omega}\| \boldsymbol{\Sigma} \mid \boldsymbol{X}$ <br> Add " $\boldsymbol{\Omega} \boldsymbol{N}$ " in the end |
| S3 | Word ends in: <br> I $\Sigma A\|I \Sigma E \Sigma\| I \Sigma E\|I \Sigma A M E\|$ <br> ILATE\|ILAN|IEANE | Remove suffix / Check exceptions / Exit | If word is: <br> IIA <br> Stem is "I $\boldsymbol{\Sigma}$ " <br> If after removal the word is: <br> ANAMIA $\mid$ A $\Theta P O\|E M \Pi A\| E \Sigma E \mid$ <br> ЕГЛКАЕ\|ЕПА|ЕАNAПА|ЕПЕ| <br> ПЕРIПA\|A@PO|इYNA@PO| <br> पANE $\mid$ K $1 E\|X A P T O П A\| E \Xi A P X A \mid$ <br> МЕТЕПЕ\|АПОКАЕ|АПЕКАЕ| <br> ЕКАЕ\|ПЕ|ПЕРІПА <br> Add "I"in the end <br> If after removal the word is: <br> AN\|AФ|ГE|ГIГANTOAФ|ГКE| <br> ДНМОКРАТ $\mid$ КОМ $\|Г К\| М\|П\|$ <br> ПОҮКАМ\|ОАО|ИАР <br> Add " $\boldsymbol{I}$ "'in the end |


| S4 | Word ends in: <br> ILS\|ILEIL|IEEI| <br> ILOYME\|ILETE|IIOYN| <br> IIOYNE | Remove suffix / Check exceptions / Exit | If after removal the word is: <br> ANAMПA\|A@PO|EMHA|ELE <br> EГЛKЛE\|EПA|ЕANAПA|EПE| <br> ПЕРIIA $\mid$ A@PO\| $\Sigma$ YNA $\Theta$ PO\| <br> पANE $\mid$ K $1 E\|X A P T O П A\| E \Xi A P X A \mid$ <br> МЕТЕПЕ\|АПОКАЕ|АПЕКАЕ| <br> ЕКАЕ\|ПЕ|ПЕРІПА <br> Add "I"in the end |
| :---: | :---: | :---: | :---: |
| S5 | Word ends in: <br> IETOE\|IETOY|IETO| <br> IETE\|ILTOI|IETRN| <br> IETOYE\|IITH|IETHE| <br> IETA\|ILTE | Remove suffix / Check exceptions / Exit | If after removal the word is: <br> $M\|\Pi\| A \Pi\|A P\| H \Delta\|K T\| \Sigma K\|\Sigma X\| Y \Psi \mid$ <br> $\Phi A\|X P\| X T\|A K T\| A O P\|A \Sigma X\| A T A \mid$ <br> AXN\|AXT|ГEM|ГYP|EMI|EYП| <br> ЕХЄ\|НФА|'НФА|КАЄ|КАК|КУИ| <br> ИYГ $\mid$ МАК $\mid$ МЕГ $\mid$ TAX $\|\Phi I \Lambda\| X \Omega P$ <br> Add " $\boldsymbol{I} \boldsymbol{\Sigma} \boldsymbol{T}$ " in the end <br> If after removal the word is: <br> पANE\| $\Sigma \mathbf{Y N A \Theta P O}\|K \Lambda E\| \Sigma E \mid$ <br> EгЛK <br> Add "I"in the end |


| S6 | Word ends in: <br> ILMO\|ILMOI|ILMOI <br> ILMOY\|ILMOYZ|ILMSN | Remove suffix / Check exceptions / Exit | * If after removal the word is: <br> AFN $\Sigma \Sigma T I K\|A T O M I K\| \Gamma N \Omega \Sigma T I K \mid$ <br> E@NIK\|EKAEKTIK| $\mathbf{\Sigma K E I T I K \| ~}$ <br> TOПIK <br> Remove "IK" from the end <br> *If after removal the word is: <br> $\Sigma E\|M E T A \Sigma E\| M I K P O \Sigma E\|E \Gamma K \wedge E\|$ <br> АПОКАЕ <br> Add " $\boldsymbol{I} \boldsymbol{\Sigma} \boldsymbol{M}$ "'in the end <br> *If after removal the word is: <br> $\triangle A N E \mid A N T I \Delta A N E$ <br> Add "I"in the end <br> *If after removal the word is: <br> AAEEANAPIN\|BYZANTIN| <br> ©EATPIN <br> Remove "IN" from the end |
| :---: | :---: | :---: | :---: |
| S7 | Word ends in: APAKI\|APAKIA|OYロAKI| OYUAKIA | Remove suffix / Check exceptions / Exit | *If after removal the word is: $X \mid \Sigma$ <br> Add "APAKI" in the end |


| S8 | Word ends in: <br> AKI\|AKIA|ITEA|ITEAE| <br> ITEEE\|ITESN|APAKI| <br> APAKIA | Remove suffix / Check exceptions / Exit | *If after removal the word is: <br> ANOP\|BAMB|BP|KAIM|KON| <br> KOP $\mid$ AABP $\mid$ ИOY $\mid$ MEP $\mid$ MOYET $\mid$ <br>  <br> ГПAN\|TZ|ФAPM|X|КАПАК| <br> AIIIФ\|AMBP|ANOP|K|ФY $\mid$ <br> КАТРАП\|КАІМ|МАА|ГАОВ $\|\Sigma \Phi\|$ <br> TEEXOEAOB <br> Add " $\boldsymbol{A K}$ " in the end <br> *If after removal the word is: <br> B\|BAA|ГIAN|ГИ||Z|HГOYMEN] <br> KAPA\|KON|MAKPYN|NYФ|| <br> ПАТЕР\|П||ЕK|TOL|TPIIOA <br> Add "ITE"in the end <br> *If after removal the word end in: KOP <br> Add "ITE"in the end |
| :---: | :---: | :---: | :---: |
| S9 | Word ends in: ILIO\|IUIA|IIIISN | Remove suffix / Check exceptions / Exit | *If after removal the word is: AISN $\mathbf{I P}\|\mathbf{O A O}\| \Psi A \Lambda$ <br> Add " $\boldsymbol{I} \boldsymbol{\Delta}$ "in the end <br> *If after removal the word iend in E\|IAIXN <br> Add " $\boldsymbol{I}$ " "in the end |
| S10 | Word ends in: <br> ILKOL\|ILKOY|ILKO <br> ILKE | Remove suffix / Check exceptions / Exit | *If after removal the word is: $\boldsymbol{\Delta}\|\boldsymbol{I B}\|$ МНN $\|\boldsymbol{P}\|$ ФРАГК $\mid$ ИYK $\mid$ OBE $\Lambda$ <br> Add " $\boldsymbol{I} \boldsymbol{\Sigma} \boldsymbol{K}$ " in the end |
| 2a | Word ends in: A4EL\|ALISN | Remove suffix / Check exceptions / Exit | If after removal the word does not end in: <br> OK\|MAM|MAN|MIAMII| ПATEP| ГIAГI|NTANT|KYP| $\mid$ EI\| ПЕЄEР <br> Add " $\boldsymbol{A}$ " " in the end |


| 2b | Word ends in： ELE $\Sigma \mid E \Delta \Omega N$ | Remove suffix／ <br> Check exceptions ／Exit | If after removal the word ends in： <br>  КРАЕП｜МІИ <br> Add＂ $\boldsymbol{E} \boldsymbol{\Delta}$＂in the end |
| :---: | :---: | :---: | :---: |
| 2c | Word ends in： OYAE $\$｜OYASN &Remove suffix／ <br> Check exceptions ／Exit&If after removal the word ends in： АРК｜КАЛІАК｜ПЕТАИ｜ИIX，｜ <br>  $X N\|\Sigma \Pi\| T P A \mid \Phi E$ <br> Add＂OY4＂in the end \hline 2d & Word ends in： E $\Omega \Sigma \mid E \Omega N$ | Remove suffix／ <br> Check exceptions ／Exit | If after removal the word is one of： АРК｜КАИІАК｜ПЕТАИ｜ИIIX｜ <br>  $X N\|\Sigma \Pi\| T P A \Gamma \mid \Phi E$ <br> Add＂$E$＂in the end |
| 3 | Word ends in： IA｜IOY｜ISN | Remove suffix | If after removal the word ends in Vowel <br> Add＂$\Gamma$＂in the end |
| 4 | Word ends in： <br> IKA｜IKO｜IKOY｜IKSN | Remove suffix／ Check exceptions ／Exit | If after removal the word ends in Vowel <br> OR <br> e of ： <br> A $\mid$｜AU｜ENA｜AMAN｜AMMOXAA｜ <br> H＠｜ANHE｜ANTIU｜ФYI｜BPSM｜ <br> ГЕР｜EЕЛА｜КААП｜КАМИIN｜ <br> KATAU｜MOYイ｜MПАN｜MПАГIAT］ <br> MНОА｜MIOЕ｜NIT｜ЕIK］ <br> ГҮNOMHИ｜ПЕTइ｜ПITE｜ <br> ПIKANT｜ПИIITTI｜ПOГTEAN｜ <br>  <br> ҮПОД｜ФIUON｜ФYイOU｜XAइ <br> Add＂ $\boldsymbol{I} \boldsymbol{K}$＂in the end |
| 5a | Word is AIAME | $\begin{aligned} & \text { Stem is } \\ & \boldsymbol{A} \boldsymbol{\Gamma A M} \\ & / \boldsymbol{E x i t} \end{aligned}$ |  |
| 5a | Word ends in ： <br> AГAME $\mid$ HEAME <br> OY氵AME｜HKAME <br> H＠HKAME | Remove suffix／ Check exceptions ／Exit |  |


| 5a | Word ends in: $A M E$ | Remove suffix / <br> Check exceptions / Exit | If after removal the word is one of: ANAП\|AПOQ|AПOK|AПOГT| ВОYВ|ЕЕЄ|OYИ|ПЕ $\mid$ \|IIKP| ПOT| $\Sigma I X \mid X$ <br> Add " $\boldsymbol{A M}$ " in the end |
| :---: | :---: | :---: | :---: |
| 5b | Word ends in: <br> ATANE\|HEANE| OYEANE |IONTANE| IOTANE|IOYNTANE| ONTANE $\mid$ OTANE $\mid$ OYNTANE\|HKANE| HEHKANE | Remove suffix / Check exceptions / Exit | If after removal the word one of: $T P \mid T \Sigma$ <br> Add " $\boldsymbol{A} \boldsymbol{\Gamma} \boldsymbol{A} \boldsymbol{N}$ " |


| 5b | Word ends in： <br> ACANE｜HEANE｜ OYEANE｜IONTANE｜ IOTANE｜IOYNTANE ONTANE｜OTANE｜ OYNTANE｜HKANE｜ H＠HKANE | Remove suffix／ Check exceptions ／Exit | If after removal the word ends in Vowel without＂$P$＂ <br> OR <br> is one of ： <br> BETEP｜BOY／K｜BPAXM $\mid \Gamma\rceil$ ДРААОYМ｜Є｜КААПОYZ｜ KAГTEイ｜КОРМОР｜ ААОПИ｜ММАМЕӨ｜М｜ МОYГOYィM｜N｜OY＾｜П｜ ПЕЛЕК｜ПИ｜ПОЛІІ｜ ПОРТОИ｜ГAPAKATE｜ EOYAT｜TEAP1AT｜OPФ｜ TEITГ｜TEOI｜ ФЛТОГTEФ｜Х｜ЧҮХОПИ｜ АГ $\mid$ ОРФ｜ГА $1 \mid$ ГЕР $\|\triangle E K\|$ पIIIИ｜AMEPIKAN｜OYP｜ ПIO｜ПOYPIT｜I｜ZQNT｜IK｜ КАЕТ｜КОП｜ИIX｜ИOY＠HP｜ MAINT｜MEイ｜$\Sigma I \Gamma\|\Sigma \Pi\|$ $\Sigma T E \Gamma\|T P A \Gamma\| T \Sigma A \Gamma\|\Phi\| E P \mid$ АААП｜АЄІГГ $\|А М Н Х\|$ ANIK $\mid$ ANOPГ $\backslash A \Pi H \Gamma\rceil$ АПIӨ｜ATEIГГ $\mid$ BA $\Sigma \mid$ ВАЕК｜ВАЄҮГA BIOMHX｜BPAXYK｜ $\mathbf{S I A T} \mid$ पIAФ｜ENOPГ $\|\Theta Y \Sigma\|$ KAПNOBIOMHX KATAГAА｜KАIB｜ KOIMAPФ｜КIB｜ МЕГГОВІОМНХ MIKPOBIOMHX $\mid$ NTAB $\mid$ ЕНРОКАІВ｜ОАІГОААМ｜ ОАОГАЯ｜ПЕNTAРФ｜ ПЕРНФ｜ПЕРITP｜ПААТ ПОАУААП｜ПОАУМНХ］ $\Sigma T E \Phi\|T A B\| T E T \mid$ ҮПЕРНФ｜ИПОКОП｜ ХАМНАО 1 АП｜ ЧНИОТАВ <br> Add＂ $\boldsymbol{A} \boldsymbol{N}$＂in the end |
| :---: | :---: | :---: | :---: |
| 5 c | Word ends in： HEETE | Remove suffix／ Check exceptions ／Exit |  |


| 5c | Word ends in: ETE | Remove suffix / <br> Check exceptions / Exit | If after removal the word ends in Vowel without " $P$ " <br> OR <br> is one of: <br> ABAP\|BEN|ENAP|ABP|AU|AQ| AN $\mid$ AПИ\|BAPON $\mid$ NTP\| $\Sigma К \mid$ KOП $\mid$ МПОР\|NIФ|ПАГ|ПАРАКАА| $\Sigma E P \Pi\|\Sigma K E \Lambda\| \Sigma Y P \Phi\|T O K\| Y\|\triangle\| E M \mid$ $\boldsymbol{\Theta A P P \|} \mid \boldsymbol{\Theta}$ <br> OR <br> ends in : <br>  EYP $\mid$ TIO\| $\boldsymbol{Y} \Pi E P \Theta\|P A \Theta\| E N \Theta \mid$ POQ\| $\Sigma \Theta\|\Pi Y P\| A I N\|\Sigma Y N \Delta\| \Sigma Y N]$ $\Sigma Y N \Theta\|X \Omega P\| \Pi O N\|B P\| K A \Theta\|E Y \Theta\|$ EKO\|NET|PON|APK|BAP|BOA| のФЕ $\boldsymbol{\Omega}$ <br> Add "ET" in the end |
| :---: | :---: | :---: | :---: |
| 5d | Word ends in: ONTA $\Sigma \mid \Omega N T A \Sigma$ | Remove suffix / Check exceptions / Exit | If after removal the word is: APX <br> add "ONT" in the end OR <br> If after removal the word is: KPE <br> add " $\boldsymbol{\Omega} \boldsymbol{N} \boldsymbol{T}$ " in the end |
| 5e | Word ends in: OMAETE\|IOMAETE | Remove suffix / Check exceptions / Exit | If after removal the word is: ON <br> add "OMAET" in the end |
| 5f | Word ends in: IELTE | Remove suffix / Check exceptions / Exit | If after removal the word is one of: $\Pi\|A \Pi\| \Sigma Y M \Pi\|A \Sigma Y M \Pi\| A K A T A \Pi \mid$ AMETAMФ <br> Add "IE $\boldsymbol{\Sigma T}$ " in the end |
| 5f | Word ends in: ELTE | Remove suffix / Check exceptions / Exit | If after removal the word is one of: AИ\|AP|EKTE $\\|Z\| M\|\Xi\|$ ПAPAKAИ $\mid$ AP\|IPO|NIL <br> Add "ELT" in the end |


| 5g | Word ends in: H@HKA\|H@HKE | H@HKE | Remove suffix / <br> Check exceptions / Exit |  |
| :---: | :---: | :---: | :---: |
| 5g | Word ends in: HKA\|HKE $\mid$ \|HKE | Remove suffix / <br> Check <br> exceptions <br> / Exit | $\square \quad$If after removal the word is <br> one of:$\boldsymbol{\operatorname { L I A Q } \boldsymbol { \Theta } \| \boldsymbol { \Theta } \| \boldsymbol { \Pi A P A K A T A \Theta } \| \boldsymbol { I P O } \boldsymbol { \Sigma } \boldsymbol { \Theta } \|}$$\boldsymbol{\Sigma Y N \Theta}$ |
| 5h | Word ends in: <br> OYェA\|OY $5 E \Sigma \mid O Y \Sigma E$ | Remove suffix / Check exceptions / Exit | If after removal the word is one of: <br> ФAPMAK\|XAA|AГK|ANAPP| ВРОМ|EKАIIT|AAMIIIU|ИEX|M| ПAT|P|I|ME $\mid$ \|ME $\Sigma A Z \mid$ YIIOTEIN\|AM|AIQ|ANHK| पEटIOZ|ENUIAФEP|UE| $\triangle E Y T E P E Y\|K A \Theta A P E Y\| \Pi \Lambda E \mid T \Sigma A$ <br> OR ends in: <br> ПОААР\|ВАЕП|ПАNТАХ|ФРУФ| MANTIL|MA』А|KYMAT $\mid$ AAX $\mid$ АНГ $\|\Phi А Г\| О М \mid П Р \Omega Т ~$ Add "OYI" in the end |


| 5 i | Word ends in： $\mathbf{A Г A}\|\mathbf{A Г E} \mathbf{\Sigma}\| \mathbf{A Г E}$ | Remove suffix／ <br> Check exceptions ／Exit | If after removal the word is one of： <br> АВАЕТ｜ПОАУФ｜АДНФ｜ПАМФ｜ P｜AГП｜AФ｜AMA $\mid$｜AMA $\Lambda \Lambda I \mid$ ANYइT $\|A \Pi E P\| A \Sigma \Pi A P\|A X A P\|$ पЕРВЕN｜ЗРОГOП｜ЕЕФ｜NEOI｜ NOMOT｜OАOI｜OMOT｜ПРОГT｜ ПРОГЛПОП｜इYMI｜$\Sigma Y N T\|T\|$ YПOT $\mid$ XAP $\mid$ AEIII $\mid$ AIMOET $\mid$ ANYП｜AПOT 1 APTIII｜UIAT｜EN｜ EПIT｜КРОКАИОП｜ГІАНРОП｜И｜ NAY｜OYイAM｜OYP｜IT｜TP｜M <br> OR <br> ends in： <br>  IP $\mid$ IOXX $\mid \Sigma M H N$ <br> BUT <br> is not one of： <br> ЧOФ｜NAYイOX <br> AND <br> does not end in： <br> KOAM <br> Add＂ $\boldsymbol{A} \boldsymbol{\Gamma}$＂in the end |
| :---: | :---: | :---: | :---: |
| 5j | Word ends in： HIE $\|\boldsymbol{H I O Y}\| \boldsymbol{H} \boldsymbol{I A}$ | Remove suffix／ Check exceptions $\square$ ／Exit | If after removal the word is one of： N｜XEPEON｜$\|\Omega \Delta E K A N\|$ EPHMON｜MEГAムON｜EПTAN <br> Add＂ $\boldsymbol{H} \boldsymbol{\Sigma}$＂in the end |
| 5k | Word ends in： HETE | Remove suffix／ <br> Check <br> exceptions <br> ／Exit | If after removal the word is one of： $A \Sigma B\|\Sigma B\| A X P\|X P\| A \Pi \Lambda \Lambda\|A E I M N\|$ GYEXP｜EYXP｜KOINOXP｜ ПААІМЧ <br> Add＂ $\boldsymbol{H} \boldsymbol{\Sigma} \boldsymbol{T}$＂in the end |
| 51 | Word ends in： OYNE｜HEOYNE｜ H＠OYNE | Remove suffix／ Check exceptions $\square$ ／Exit | If after removal the word is one of： N｜P｜इIII｜$\Sigma$ TPABOMOYTE｜ KAKOMOYTE｜EESN <br> Add＂OYN＂in the end |


| 51 | Word ends in: OYME\|HEOYME| H@OYME | Remove suffix / <br> Check exceptions $\square$ / Exit | If after removal the word is one of: ПAPAEOY $\Sigma\|\Phi\| X\|\Omega P I O \Pi \Lambda\| A Z \mid$ AMAOEOYE\|AEOYE <br> Add "OYM" in the end |  |
| :---: | :---: | :---: | :---: | :---: |
| 6 | Word ends in: <br> MATA\|MATRN|MATOE | Remove suffix | Always <br> Add " $\boldsymbol{M} \boldsymbol{A}$ " in the end |  |
| 6 | Word ends in: <br> A\|ACATE $\mid$ AГAN $\mid$ AEI $\mid$ AMAI $\|A N\| A \Sigma\|A \Sigma A I\| A T A I\|A \Omega\| E\|E I\| E I \Sigma \mid$ EITE\|E $\mathbf{E A I \| E \Sigma \| E T A I \| \| I \| I E M A I \| I E M A \Sigma T E \| I E T A I \| I E \Sigma A I \| ~}$ IE EAETE\|IOMAETAN|IOMOYN|IOMOYNA|IONTAN] IONTOYEAN|IOEAETAN|IOEAETE|IOEOYN|IOEOYNA| IOTAN|IOYMA|IOYMAETE|IOYNTAI|IOYNTAN|H|HDEL| HUSN|H@EI|H@EIL|H@EITE|H@HKATE|H@HKAN|H@OYN| H@S|HKATE|HKAN|H $\Sigma\|H \Sigma A N\| H \Sigma A T E\|H \Sigma E I\| H \Sigma E \Sigma\|H \Sigma O Y N\|$ H $\Sigma \Omega\|O\| O I\|O M A I\| O M A \Sigma T A N\|O M O Y N\| O M O Y N A\|O N T A I\|$ ONTAN\|ONTOYEAN $\|O \Sigma\| O \Sigma A \Sigma T A N\|O \Sigma A \Sigma T E\| O \Sigma O Y N \mid$ OLOYNA\|OTAN|OY|OYMAI|OYMAETE|OYN|OYNTAI| OYNTAN|OYE|OYEAN|OYEATE|Y|YE| $\mid$ \| $\Omega N$ |  |  | Remove suffix |
| 7 | Word ends in: <br> ELTEP\|EETAT $\mid$ OTEP $\mid$ <br> OTAT\|YTEP|YTAT|ITEP| <br> תTAT | Remove suffix |  |  |
| UL2 | Word has converted letters to upper case from step UL1 | Convert these letters back to lower case |  |  |

## 6. Final Evaluation

After the implementation of all the modifications that would improve the existing algorithm, we begun to evaluate our revised algorithm. We used the same word list and the same applications we created for testing the initial algorithm. By doing so, we made sure that the statistics produced can be directly comparable.

As a result of our modifications, we ended with a stronger stemmer. Although the two stemmers leave unchanged roughly the same number of words, our modified version produces fewer and bigger conflation classes by altering more letters in every word, on average. Table 9 presents the statistics gathered after executing both algorithms against a list of 574.621 Greek words.

Table 9: Statistics: Comparison of the original and revised algorithm

|  | Original Stemmer by <br> Ntais | Our modified <br> Stemmer |
| :--- | :--- | :--- |
| Mean number of words per <br> conflation class | 4,055 | 5,664 |
| Index compression factor | $75,34 \%$ | 82,34 |
| Ratio of unchanged to total words | $2 \%$ | $2 \%$ |
| Mean modified Humming Distance | 2.441 | 2,916 |
| Median Modified Humming Distance | 2 | 2 |
| Correct Stems (for a sample of 12468 <br> words) | $10885(87,3 \%)$ | $11669(93,52 \%)$ |
|  | Distribution of Stemming errors per algorithm |  |
|  | Original Stemmer by | Our modified |
| Ntais | Stemmer |  |
| Understemming Errors <br> (Section 3.2) | $88,44 \%$ | $23,67 \%$ |
| Overstemming Errors <br> (Section 3.2) | $11,56 \%$ | $76,33 \%$ |
| Number of different stems generated <br> by the two stemmers (for a sample of <br> 574.62 words) | $35885(6,24 \%)$ |  |

It is worth noting that, according to our tests, the majority of the errors of the initial algorithm had to do with understemming ( $88,44 \%$ ). Our algorithm produces more overstemming errors $(76,33 \%)$ despite of the fact that the total number of errors is reduced. The two stemmers produced 35885 different stems for the same list of words. In addition the number of executions steps was increased from 29 in the original algorithm to 42 . Ten of the new execution steps have to do with the 72 new stems that were added while the remaining three deal with stop-word removal and lower to upper and upper to lower case treatment.

Although the number of steps was increased by approximately $44 \%$, the algorithm now executes 23.17 steps on average. The reason is that while the original algorithm always executes all of its 29 steps, our modified algorithms returns the correct stem and then exits earlier if the remaining rules are not going to modify the word any further. For example, if a rule that treats verb suffixes is evaluated and executed, it is certain that no rules that deal with noun suffixes will be executed later. This modification will certainly offer better running times in any implementation of our algorithm. Nevertheless, we avoid mentioning any running times and we only refer to the average number of steps executed. The reason is that the actual running time of an algorithm depends on factors like the implementation language, coding style and coding efficiency, hardware, current load of the system and operating system among others. Because of that, we believe that comparing running times of our tests between our algorithm and the initial algorithm by Ntais would be misleading.

## 7. Conclusion

Our purpose has been to evaluate and improve the existing stemming algorithm of the Greek language. After an evaluation of the results that the original algorithm provided, we incrementally improved it by adding new rules and exceptions. Overall, we managed to gain significant improvements in completeness and accuracy to the existing algorithm.

After manually checking the output of the stemmer we conclude that our new algorithm returns more correct results than its predecessor. Due to our efforts the understemming and overstemming errors are less. The new algorithm is more complete since it supports most tenses and correctly stems suffixes that were not included before, like diminutives and others.

Moreover, we offered a more usable implementation that can be used with slight modifications or even directly. Due to the implementation language that we chose, PHP, our implementation can be used by any web or non web application that may require stemming of Greek words. Our implementation is directly usable by any kind of application, web or not, linguistic or IR related that requires stemming for Greek.

To answer our research question, we conclude that the addition of more suffixes is attainable but the effort required for each additional suffix increases geometrically. The previous algorithm already deals with the majority of suffixes that can be found in the Greek grammar. Since the majority of cases is already covered, each addition to the stemmer requires a considerable amount of evaluating the stemmer's output and searching in dictionaries for possible exceptions.

Despite of the fact that we are pleased with the outcome of our efforts, there is a lot of space for improvements. Our algorithm, like its predecessor, is not dealing with the moving tone-mark issue. Any attempt to deal with issue will probably require a considerable amount of effort due to the complexity of the Greek language. Although it will not provide any great improvement in precision, a stemmer must successfully deal with this issue in order to be considered complete. Furthermore, the stemmer can be enhanced by the addition of more suffixes and exceptions. In addition to the 158 suffixes of the initial algorithm, we added rules for 72 more. Although the majority of the cases is covered, more rules can be added in order to produce a more complete stemmer. Finally, the stemmer can be more thoroughly tested. Although we used metrics
to measure our performance and compare it with its predecessor, we had to manually check the stems generated for overstemming and understemming errors. This is a task that can only be done manually, by experts of the Greek language. Although we had an enormous collection of Greek words, we were able to test only one small portion of it, as we mention in Table 9. A more complete testing will indicate more errors and can serve as a basis for further improvement.

We believe that this thesis work contributed in stemming research by continuing and extending the work done by others in the past and by offering a stemmer for the Greek language which others can use and extend even more.

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## Appendices

## Appendix A: Verb Conjugation Classes In Greek

## Verbs of $1^{\text {st }}$ Conjugation Classes

| The verb" "to read" in Greek Active Voice / Past Tenses |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  <br> (Past Perfect) |  | Aópıбтos (Simple Past) |  | Парататько́я <br> (Past continuous ) |  |
| Singular | عíq $\alpha$ <br> $\delta 1 \alpha \beta \dot{\alpha} \sigma \varepsilon!$ | (I) had studied | $\delta \iota \alpha$ ¢ $\beta \boldsymbol{\sigma} \alpha$ | (I) studied | $\delta$ ı $\alpha$ 人 $\alpha$ ¢ $\alpha$ | (I) was studying |
|  | ع́хє६ऽ <br> $\delta 1 \alpha \beta \dot{\alpha} \sigma \varepsilon!$ | (you) had studied | $\delta 1$ 人́ß $\alpha \sigma \varepsilon \varsigma$ | (you) studied | $\delta 1 \alpha{ }^{1} \beta \alpha \zeta \varepsilon ¢$ | (you) were studying |
|  | ві́xє <br> $\delta \iota \alpha \beta \dot{\sigma} \sigma \varepsilon \iota$ | (he) had studied |  | (he) studied | $\delta \dot{\alpha} \beta \alpha \zeta ¢$ | (he) was studying |
| Plural | عі́ха $\mu$ <br> $\delta 1 \alpha \beta \alpha ́ \sigma \varepsilon!$ | (we) had studied | $\delta i \alpha \beta \alpha \alpha^{\prime} \alpha \mu \varepsilon$ | (we) studied | $\delta 1 \alpha \beta \alpha{ }^{\text {a }}$ ¢ $\alpha \mu \varepsilon$ | (we) were studying |
|  | عí $\alpha \tau \varepsilon$ <br> $\delta 1 \alpha \beta \dot{\alpha} \sigma \varepsilon!$ | (you) had studied | $\delta i \alpha \beta \alpha \alpha^{\prime} \alpha \tau \varepsilon$ | (you) <br> studied | $\delta 1 \alpha \beta \alpha{ }^{\text {a }}$ ¢ $\alpha \tau$ | (you) were studying |
|  | sí $\alpha \nu$ <br> $\delta 1 \alpha \beta \alpha ́ \sigma \varepsilon t$ | (they) had studied | $\delta 1 \alpha \beta \alpha ́ \sigma \alpha v \varepsilon$ <br> ( $\delta$ tó $\beta \alpha \boldsymbol{\sigma} \alpha v$ ) | (they) <br> studied |  | (they) were studying |


| The verb "to read" in Greek Active Voice / Present Tenses |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
|  | Eveбтळ́tas <br> (Present) |  | Параквíдєvоя <br> (Present Perfect) |  |
| Singular | $\delta ı \alpha \beta \dot{\zeta} \boldsymbol{\omega}$ | (I) read /am reading | $\dot{\varepsilon} \chi \omega$ ¢ı $\alpha \beta \alpha \dot{\sigma} \boldsymbol{\varepsilon}$ ı | (I) have read /have been reading |
|  | $\delta 1 \alpha \beta \alpha ́ \zeta \varepsilon ı \zeta$ | (you) read / are reading |  | (you) have read/have been reading |
|  | $\delta ı \alpha \beta \alpha{ }^{\text {a }}$ ¢ $¢$ | (he) reads/ is reading | $\dot{\varepsilon} \chi \varepsilon 1 \delta 1 \alpha \beta \alpha \dot{\sigma} \varepsilon$ ו | (he) has read/has been reading |
| Plural | $\delta i \alpha \beta \dot{\zeta} \mathbf{0} \mathbf{0} \mu \boldsymbol{\varepsilon}$ | (we) read/ are reading |  | (we) have read/have been reading |
|  | $\delta 1 \alpha \beta \alpha ́ \zeta \varepsilon \tau \varepsilon$ | (you) read / are reading | غ́ $\chi \varepsilon \tau \varepsilon \delta\langle\alpha \beta \dot{\alpha} \boldsymbol{\sigma} \varepsilon$ ı | (you) have read/have been reading |
|  | $\delta$ ¢ $\alpha \beta \alpha \zeta^{\text {covv }}$ | (they) read /are reading |  | (they) have read/have been reading |


| The verb" "to read" in Greek Active Voice / Future Tenses |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Еگккодоvөŋтıко́я <br> Mé $\lambda \lambda o v \tau \alpha \varsigma$ <br> ( Future Continuous) |  |  (Simple Future) |  | $\Sigma \nu v \tau \varepsilon \lambda \varepsilon \sigma \mu \varepsilon ́ v o \varsigma ~ M \varepsilon ́ \lambda \lambda о v \tau \alpha \varsigma$ <br> (Future Perfect/Imperfect) |  |
| Singular | $\theta \alpha \delta 1 \alpha \beta \alpha \zeta^{\prime} \omega$ | (I) will be reading | $\theta \alpha \delta 1 \alpha \beta \alpha \alpha^{\prime} \omega$ | (I) will read | $\theta \alpha \varepsilon ́ \chi \omega$ <br> $\delta 1 \alpha \beta \dot{\alpha} \sigma \varepsilon \iota$ | (I) will have read / have been reading |
|  | $\theta \alpha \delta 1 \alpha \beta \dot{\zeta} \zeta \boldsymbol{\varepsilon}$ ¢¢ | (you) will be reading | $\theta \alpha$ ठı $\alpha \beta \dot{\alpha} \boldsymbol{\sigma}$ ¢ı | (you) will read | $\theta \alpha$ غ́ $\chi \varepsilon ા \varsigma$ <br> $\delta 1 \alpha \beta \dot{\alpha} \sigma \varepsilon \iota$ | (you) will have read/have been reading |
|  |  | (he) will be reading | $\theta \alpha \delta 1 \alpha \beta \alpha \alpha^{\prime} \boldsymbol{\varepsilon}$ ı | (he) will read | $\theta \alpha$ ह́ $\chi \varepsilon 1$ $\delta 1 \alpha \beta \dot{\alpha} \sigma \varepsilon \iota$ | (he) will have read have been reading |
| Plural | $\theta \alpha$ $\delta ı \beta \alpha \dot{\zeta} \zeta \boldsymbol{0} \boldsymbol{\mu} \varepsilon$ | (we )will reading | $\theta \alpha$ $\delta 1 \alpha \beta \alpha ́ \sigma 0 v \mu \varepsilon$ | (we) will read | $\theta \alpha$ غ́ $\chi о \cup \mu \varepsilon$ $\delta 1 \alpha \beta \dot{\alpha} \sigma \varepsilon!$ | (we) will have read/have been reading |
|  | $\theta \alpha$ $\delta 1 \alpha \beta \alpha \alpha_{\boldsymbol{\varepsilon}} \boldsymbol{\tau} \varepsilon$ | (you)will be reading | $\theta \alpha \delta 1 \alpha \beta \alpha \sigma^{\boldsymbol{\sigma} \boldsymbol{\varepsilon} \boldsymbol{\tau} \varepsilon}$ | (you) will read | $\theta \alpha \varepsilon ́ \chi \varepsilon \tau \varepsilon$ <br> $\delta 1 \alpha \beta \alpha ́ \sigma \varepsilon t$ | (you) will have read/have been reading |
|  | $\theta \alpha$ $\delta 1 \alpha \beta \alpha ́ \zeta$ ovv | (they) will be reading | $\theta \alpha$ $\delta 1 \alpha \beta \alpha ́ \sigma 0 v v \varepsilon$ | (they) will read | $\theta \alpha$ غ́zouv $\delta 1 \alpha \beta \alpha ́ \sigma \varepsilon t$ | (they) will have read/have been reading |


| The verb "to read" in Greek Passive Voice / Past Tenses |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  <br> (Past Perfect) |  | Aópıotoç (Simple Past) |  | Парататько́я (Past Continuous ) |  |
| Singular | عíz $\alpha$ <br> $\delta 1 \alpha \beta \alpha \sigma \tau \varepsilon$ í | (I) had been studied | $\delta 1 \alpha \beta \alpha ́ \sigma \tau \eta \kappa \alpha$ | (I) was studied | $\delta\langle\alpha \beta$ ¢óóuovv | (I) was being studied |
|  | عí $\chi \varepsilon \varsigma$ <br> $\delta 1 \alpha \beta \alpha \sigma \tau \varepsilon$ í | (you) had been studied | $\delta 1 \alpha \beta \alpha ́ \sigma \tau \eta \kappa \varepsilon \varsigma$ | (you) were studied | $\delta 1 \alpha \beta \alpha$ ̧óбovv | (you) were being studied |
|  | દíxє <br> $\delta 1 \alpha \beta \alpha \sigma \tau \varepsilon$ í | (he) had been studied | $\delta 1 \alpha \beta \dot{\alpha} \sigma \tau \eta \kappa \varepsilon$ | (he) was studied | $\delta 1 \alpha \beta \alpha$ ¢ó $\boldsymbol{\tau} \boldsymbol{\nu}$ | (he) was being studied |
| Plural | ві́х $\alpha \mu \varepsilon$ <br> $\delta i \alpha \beta \alpha \sigma \tau \varepsilon$ í | (we) had been studied | $\delta 1 \alpha \beta \alpha \sigma \tau \eta \prime \kappa \alpha \mu \varepsilon$ | (we) were studied | $\delta 1 \alpha \beta \alpha$ ¢ó $\mu \boldsymbol{\alpha} \boldsymbol{\tau} \boldsymbol{\tau}$ | (we) were being studied |
|  | દі́ $\alpha \boldsymbol{\alpha} \varepsilon$ <br> $\delta 1 \alpha \beta \alpha \sigma \tau \varepsilon$ í | (you) had been studied |  | (you) were studied | $\delta 1 \alpha \beta \alpha \zeta$ ¢́ $\sigma \alpha \sigma \tau \varepsilon$ | (you) were being studied |
|  | عí $\alpha v$ <br> $\delta 1 \alpha \beta \alpha \sigma \tau \varepsilon$ í | (they) had been studied | $\delta 1 \alpha \beta \alpha \sigma \tau \eta ́ \kappa \alpha v \varepsilon$ <br> ( $\delta 1 \alpha \beta \alpha ́ \sigma \tau \eta \kappa \alpha v)$ | (they) <br> were <br> studied |  | (they) were being studied |


| The verb "to read" in Greek Active Voice / Present Tenses |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
|  | Evєбтळ́тая <br> (Present) |  | Паракє́́иєvоя <br> (Present Perfect) |  |
| Singular | бıаßá̧ouaı | (I) am read /am being read |  | (I) have been read |
|  | ঠıßßá̧ıбоı | (you) are read / are being read |  | (you) have been read |
|  |  | (he) is read/ is being read |  | (he) has been read |
| Plural |  | (we) are read/ are being read | غ́хоบนє $\delta<\alpha \beta \alpha \sigma \tau \varepsilon$ í | (we) have been read |
|  | $\delta ı \beta \alpha \dot{\zeta} \varepsilon \sigma \tau \varepsilon$ | (you) are read / are being read | $\dot{\varepsilon} \chi \varepsilon \tau \varepsilon \delta \delta \alpha \beta \alpha \sigma \tau \varepsilon i$ | (you) have been read |
|  | ঠıaßágovtaı | (they) are read /are being read | غ́xouv $\delta ı \alpha \beta \alpha \sigma \tau \varepsilon i ́$ | (they) have been read |


| The verb "to read" in Greek Active Voice / Future Tenses |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  <br> Méдддovtas <br> ( Future Continuous) |  | $\Sigma \tau \tau \gamma \mu \iota \alpha i ́ o c ̧ ~ М е ́ \lambda \lambda о v \tau \alpha \varsigma ~$ <br> (Simple Future) |  | $\Sigma \nu v \tau \varepsilon \lambda \varepsilon \sigma \mu \varepsilon ́ v o \varsigma ~ M \varepsilon ́ \lambda \lambda д o v \tau \alpha \varsigma$ <br> (Future Perfect/Imperfect) |  |
| Singular |  | (I) will be being reading | $\theta \alpha \delta 1 \alpha \beta \alpha \sigma \tau \omega \prime$ | (I) will be read | $\theta \alpha \varepsilon ́ \chi \omega$ <br> $\delta 1 \alpha \beta \alpha \sigma \tau \varepsilon$ í | (I) will have been read |
|  | $\theta \alpha \delta ı \alpha \beta \alpha{ }^{\text {a }}$ ¢ $\boldsymbol{\varepsilon} \boldsymbol{\alpha} \boldsymbol{l}$ | (you) will be being read | $\theta \alpha$ $\delta \alpha \beta \alpha \sigma \tau \varepsilon$ ќs | (you) will be read | $\theta \alpha$ غ́ $\chi \varepsilon 1 \varsigma$ <br> $\delta 1 \alpha \beta \alpha \sigma \tau \varepsilon$ í | (you) will have been read |
|  | $\theta \alpha \delta ı \alpha \beta \alpha{ }^{\text {a }}$ ¢ $¢ \tau \alpha \mathfrak{l}$ | (he) will be being read | $\theta \alpha \delta 1 \alpha \beta \alpha \sigma \tau \varepsilon$ í | (he) will be read | $\theta \alpha$ غ́ $\chi \varepsilon \iota$ <br> $\delta 1 \alpha \beta \alpha \sigma \tau \varepsilon$ í | (he) will have been read |
| Plural | $\theta \alpha$ $\delta 1 \alpha \beta \alpha \zeta{ }^{\circ} \boldsymbol{\mu} \boldsymbol{\mu} \alpha \sigma \tau \varepsilon$ | (we) will be being read | $\theta \alpha$ $\delta 1 \alpha \beta \alpha \sigma \tau 0 v ์ \mu \varepsilon$ | (we) will be read | $\theta \alpha$ غ́ $\chi \circ \cup \mu \varepsilon$ $\delta 1 \alpha \beta \alpha \sigma \tau \varepsilon$ í | (we) will have been read |
|  | $\theta \alpha$ <br>  | (you) will be being read | $\theta \alpha$ <br> $\delta 1 \alpha \beta \alpha \sigma \tau \varepsilon i ́ \tau \varepsilon$ | (you) will be read | $\theta \alpha$ غ́ $\chi \varepsilon \tau \varepsilon$ <br> $\delta 1 \alpha \beta \alpha \sigma \tau \varepsilon$ í | (you) will have been read |
|  | $\theta \alpha$ $\delta 1 \alpha \beta \alpha ́ \zeta 0 v \tau \alpha \iota$ | (they) will be being read | $\theta \alpha$ <br> $\delta 1 \alpha \beta \alpha \sigma \tau 0 v ์ v$ | (they) will be read | $\theta \alpha$ غ́ $\chi o v v$ <br> $\delta 1 \alpha \beta \alpha \sigma \tau \varepsilon$ í | (they) will have been read |

## Verbs of $2^{\text {nd }}$ Conjugation Class

| The verb "to love" in Greek Active Voice / Past Tenses |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  (Past Perfect) |  | Aópıбтos (Simple <br> Past) |  | Парататıко́я <br> (Past continuous ) |  |
| Singular | عí $\alpha$ <br> $\alpha \gamma \alpha \pi \eta ́ \sigma \varepsilon \iota$ | (I) had loved | $\alpha \gamma \alpha \dot{\pi} \eta \boldsymbol{\sigma} \boldsymbol{\alpha}$ | (I) loved | $\alpha \gamma \alpha \pi о v ́ \sigma \alpha$ <br> ( $\alpha \gamma \dot{\alpha} \pi \alpha \gamma \boldsymbol{\alpha}$ ) | (I) was loving |
|  | عíxยऽ <br> $\alpha \gamma \alpha \pi \eta{ }^{\sigma} \sigma \varepsilon!$ | (you) had loved | $\alpha \gamma \alpha ́ \pi \eta \sigma \varepsilon \varsigma$ | (you) loved | $\alpha \gamma \alpha \pi о$ б́бєऽ <br> ( $\alpha \gamma \alpha \dot{\alpha} \alpha \gamma \varepsilon \varsigma)$ | (you) were loving |
|  | عíx <br> $\alpha \gamma \alpha \pi \eta ́ \sigma \varepsilon \iota$ | (he) had loved | $\alpha \gamma \alpha \chi^{\prime} \boldsymbol{\eta} \boldsymbol{\sigma} \boldsymbol{\varepsilon}$ | (he) loved | $\alpha \gamma \alpha \pi$ ои́ $\sigma \varepsilon$ <br> ( $\alpha \gamma \alpha \dot{\alpha} \pi \alpha \gamma \varepsilon)$ | (he) was loving |
| Plural | вíq$\alpha \mu \varepsilon$ <br> $\alpha \gamma \alpha \pi \eta ́ \sigma \varepsilon!$ | (we) had loved | $\alpha \gamma \alpha \pi \eta ์ \sigma \alpha \mu \varepsilon$ | (we) loved | $\alpha \gamma \alpha \pi о$ $\sigma \alpha \mu \varepsilon$ <br> ( $\alpha \gamma \alpha \pi \alpha ́ \gamma \alpha \mu \varepsilon)$ | (we) were loving |
|  | $\varepsilon$ 白 $\alpha \tau \varepsilon$ <br> $\alpha \gamma \alpha \pi \eta ́ \sigma \varepsilon \iota$ | (you) had loved | $\alpha \gamma \alpha \pi \eta ์ \sigma \alpha \tau \varepsilon$ | (you) loved | $\alpha \gamma \alpha \pi о$ б́ $\alpha \tau \varepsilon$ <br> ( $\alpha \gamma \alpha \pi \alpha ́ \gamma \alpha \tau \varepsilon)$ | (you) were loving |
|  | عí $\alpha v$ <br> $\alpha \gamma \alpha \pi \eta ́ \sigma \varepsilon \iota$ | (they) had loved | $\alpha \gamma \alpha{ }^{\prime} \pi \eta \sigma \alpha \nu$ | (they) loved | $\alpha \gamma \alpha \pi о v ́ \sigma \alpha v$ <br> ( $\alpha \gamma \alpha \dot{\alpha} \pi \alpha \gamma \alpha v)$ | (they) were loving |

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| The verb "to love" in Greek Active Voice / Present Tenses |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
|  | Eveбтळ́тas <br> (Present) |  | Паракє́́ $\varepsilon$ гоя <br> (Present Perfect) |  |
| Singular | $\alpha \gamma \alpha \pi \dot{\omega}$ <br> ( $\alpha \gamma \alpha \pi \alpha ́ \omega)$ | (I) love /am loving |  | (I) have loved /have been loving |
|  | $\alpha \gamma \alpha \pi \alpha ́ \varsigma$ | (you) love / are loving |  | (you) have loved/have been loving |
|  | $\alpha \gamma \alpha \pi \boldsymbol{\alpha}$ <br> ( $\alpha \gamma \alpha \pi \alpha ́ \varepsilon$ ) | (he) loves/ is loving |  | (he) has loved/has been loving |
| Plural | $\alpha \gamma \alpha \pi$ ои́ $\boldsymbol{\mu} \varepsilon$ <br> ( $\alpha \gamma \alpha \pi \alpha ́ \mu \varepsilon)$ | (we) love / are loving | $\dot{\varepsilon} \chi<\cup \mu \varepsilon \alpha \gamma \alpha \pi \eta ์ \sigma \varepsilon!$ | (we) have loved/have been loving |
|  | $\alpha \gamma \alpha \pi \alpha ́ \tau \varepsilon$ | (you) love / are loving | $\varepsilon ́ \chi \varepsilon \tau \varepsilon \alpha \gamma \alpha \pi \eta ์ \sigma \varepsilon!$ | (you) have loved/have been loving |
|  | $\alpha \gamma \alpha \pi \alpha ́ v \varepsilon$ | (they) love /are loving |  | (they) have <br> loved/have been <br> loving |


| The verb" "to love" in Greek Active Voice / Future Tenses |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Еگакодоv日ŋтıко́я <br> Mé $\lambda \lambda o v \tau \alpha \varsigma$ <br> ( Future Continuous) |  |  (Simple Future) |  |  <br> (Future Perfect/Imperfect) |  |
| Singular | $\theta \alpha \alpha \gamma \alpha \pi \omega$ <br> ( $\theta \alpha \alpha \gamma \alpha \pi \alpha ́ \omega)$ | (I) will be loving | $\theta \alpha \alpha \gamma \alpha \pi \eta ์ \sigma \omega$ | (I) will love | $\theta \alpha \varepsilon ́ \chi \omega$ $\alpha \gamma \alpha \pi \eta ์ \sigma \varepsilon \iota$ | (I) will have loved /have been loving |
|  | $\theta \alpha \alpha \gamma \alpha \pi \alpha ́ \varsigma$ | (you) will be loving | $\theta \alpha \alpha \gamma \alpha \pi \eta ์ \sigma \varepsilon เ \varsigma$ | (you) will love | $\theta \alpha$ ह́ $\chi \varepsilon \iota$ $\alpha \gamma \alpha \pi \eta ́ \sigma \varepsilon \iota$ | (you) will have loved/have been loving |
|  | $\theta \alpha \alpha \gamma \alpha \pi \alpha ́ \varepsilon \backslash$ | (he) will be loving | $\theta \alpha \alpha \gamma \alpha \pi \underline{\sigma} \boldsymbol{\varepsilon} \boldsymbol{\varepsilon}$ ı | (he) will love | $\theta \alpha$ غ́ $\chi \varepsilon ı$ $\alpha \gamma \alpha \pi \eta ́ \sigma \varepsilon \iota$ | (he) will has loved/has been loving |
| Plural | $\theta \alpha \alpha \gamma \alpha \pi \boldsymbol{v} \mu \varepsilon$ ( $\theta \alpha \alpha \gamma \alpha \pi \alpha ́ \mu \varepsilon)$ | (we )will <br> loving | $\theta \alpha$ $\alpha \gamma \alpha \pi \eta ́ \sigma о \nu \mu \varepsilon$ | (we) will love | $\theta \alpha$ દ́ $\chi о ч \mu \varepsilon$ $\alpha \gamma \alpha \pi \eta(\sigma \varepsilon \iota$ | (we) will have loved/have been loving |
|  | $\theta \alpha \alpha \gamma \alpha \pi \boldsymbol{\alpha} \boldsymbol{\tau} \boldsymbol{\varepsilon}$ | (you)will be loving | $\theta \alpha$ <br> $\alpha \gamma \alpha \pi \eta ́ \sigma \varepsilon \tau \varepsilon$ | (you) will love | $\theta \alpha$ غ́ $\chi \varepsilon \tau \varepsilon$ $\alpha \gamma \alpha \pi \eta ́ \sigma \varepsilon \iota$ | (you) will have loved/have been loving |
|  | $\theta \alpha \alpha \gamma \alpha \pi \boldsymbol{v} v$ ( $\theta \alpha \alpha \gamma \alpha \pi \alpha ́ v \varepsilon)$ | (they) will be loving | $\theta \alpha \alpha \gamma \alpha \pi \eta ์ \sigma 0 v v$ | (they) will love | $\theta \alpha$ é $\chi o u v$ $\alpha \gamma \alpha \pi \eta ́ \sigma \varepsilon \iota$ | (they) will have loved/have been loving |


| The verb" $t \mathrm{to}$ love" in Greek Passive Voice / Past Tenses |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  (Past Perfect) |  | Aópıбtos (Simple <br> Past) |  | Парататıко́s <br> (Past continuous ) |  |
| Singular | عíx $\alpha \gamma \alpha \pi \eta \theta \varepsilon$ í | (I) had been loved | $\alpha \gamma \alpha \pi \grave{\dagger} \theta \eta \kappa \alpha$ | (I) was <br> loved | aүaлıónovv | (I) was being loved |
|  | عíx६ऽ <br> $\alpha \gamma \alpha \pi \eta \theta z i ́$ | (you) had been loved | $\alpha \gamma \alpha \pi \mathfrak{\eta} \theta \eta \kappa \varepsilon \varsigma$ | (you) were <br> loved | aүaлıócovv | (you) were being loved |
|  | عíx $\alpha \gamma \alpha \pi \eta \theta \varepsilon i ́$ | (he) had been loved | $\alpha \gamma \alpha \pi \dot{\eta}$ өضк | (he) was loved | $\alpha \gamma \alpha \pi$ เótov | (he) was <br> being loved |
| Plural | ві́ $\alpha \mu \varepsilon$ <br> $\alpha \gamma \alpha \pi \eta \theta \varepsilon i ́$ | (we) had <br> been loved | $\alpha \gamma \alpha \pi \eta$ ө́¢канц | (we) were loved | $\alpha \gamma \alpha \pi \iota o ́ \mu \alpha \sigma \tau \varepsilon$ <br> ( $\alpha \gamma \alpha \pi \iota o ́ \mu \alpha \sigma \tau \alpha v)$ | (we) were <br> being loved |
|  | عí $\alpha \tau \varepsilon$ <br> $\alpha \gamma \alpha \pi \eta \theta \varepsilon i$ | (you) had been loved |  | (you) were <br> loved | $\alpha \gamma \alpha \pi เ$ о́ $\sigma \sigma \tau \varepsilon$ ( $\alpha \gamma \alpha \pi \iota o ́ \sigma \alpha \sigma \tau \alpha v)$ | (you) were <br> being loved |
|  | عí $\alpha v$ <br> $\alpha \gamma \alpha \pi \eta \theta \varepsilon$ í | (they) had been loved |  | (they) were <br> loved | $\alpha \gamma \alpha \pi$ ıóvtav <br> ( $\alpha \gamma \alpha \pi$ ıóvтova $\alpha v$ ) | (they) were <br> being loved |


| The verb "to love" in Greek <br> Passive Voice / Present Tenses |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
|  | Eveбтळ́та¢ (Present) |  | Парокє́́иєvоя <br> (Present Perfect) |  |
| Singular | $\alpha \gamma \alpha \pi ı$ ¢́ $\mu \boldsymbol{\alpha}$ | (I) am being loved | $\varepsilon ̇ \chi \omega \alpha \gamma \alpha \pi \eta \theta \varepsilon \varepsilon ́$ | (I) have been loved |
|  |  | (you) are being loved | غ́ $\chi \varepsilon 1 \zeta \alpha \gamma \alpha \pi \eta \theta \varepsilon$ cí | (you) have been loved |
|  | $\alpha \gamma \alpha \pi \iota$ ı́̇ $\alpha$ ı | (he) is being loved | غ́ $\chi \varepsilon 1 \sim \gamma \alpha \pi \eta \theta \varepsilon$ ı́ | (he) has been loved |
| Plural | $\alpha \gamma \alpha \pi \iota o ́ \mu \alpha \sigma \tau \varepsilon$ | (we) are being loved | غ́ $\chi 0 \cup \mu \varepsilon \alpha \gamma \alpha \pi \eta \theta \varepsilon$ ć | (we) have been loved |
|  | $\alpha \gamma \alpha \pi \iota$ ó $\sigma \alpha \sigma \tau \varepsilon$ | (you) are being loved | $\varepsilon$ é $\chi \varepsilon \tau \varepsilon \alpha \gamma \alpha \pi \eta \theta \varepsilon i ́$ | (you) have been loved |
|  | $\alpha \gamma \alpha \pi ı 0 v ์ v \tau \alpha \iota$ | (they) are being loved | غ́zovv $\alpha \gamma \alpha \pi \eta$ Өcí | (they) have been loved |


| The verb "to love" in Greek <br> Passive Voice / Future Tenses |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  <br> Mé $\lambda \lambda \lambda^{\prime} \tau \tau \alpha \varsigma$ <br> ( Future Continuous) |  | $\Sigma \tau ı \gamma \mu \iota \alpha i ́ o g ~ M \check{́ c \lambda} \lambda о v \tau \alpha \varsigma$ (Simple Future) |  | इvviє $\lambda \varepsilon \sigma \mu \varepsilon ́ v o s$ <br> Mé $\lambda \lambda 0 v \tau \alpha \varsigma$ <br> (Future <br> Perfect/Imperfect) |  |
| Singular | $\theta \alpha \alpha \gamma \alpha \pi \mathrm{t}$ ¢́ $\mu \boldsymbol{\alpha}$ | (I) will be loved | $\theta \alpha \alpha \gamma \alpha \pi \eta \boldsymbol{\theta} \boldsymbol{\omega}$ | (I) will love | $\theta \alpha \varepsilon ́ \chi \omega$ $\alpha \gamma \alpha \pi \eta \theta \varepsilon$ í | (I) will have been loved |
|  | $\theta \alpha \alpha \gamma \alpha \pi \mathrm{t}$ ¢́б $\alpha \boldsymbol{l}$ | (you) will be loved | $\theta \alpha \alpha \gamma \alpha \pi \eta \theta \varepsilon \varepsilon^{\prime} \varsigma$ | (you) will love | $\theta \alpha$ ह́ $\chi \varepsilon 1 \varsigma$ $\alpha \gamma \alpha \pi \eta \theta \varepsilon i ́$ | (you) will have been loved |
|  | $\theta \alpha \alpha \gamma \alpha \pi เ$ ¢́ $\boldsymbol{\alpha} \boldsymbol{\chi}$ | (he) will be loved | $\theta \alpha \alpha \gamma \alpha \pi \eta \theta \varepsilon$ í | (he) will love | $\theta \alpha \varepsilon ́ \chi \varepsilon 1$ $\alpha \gamma \alpha \pi \eta \theta \varepsilon i ́$ | (he) will have been loved |
| Plural | $\theta \alpha$ $\alpha \gamma \alpha \pi$ ıó $\mu \alpha \sigma \tau \varepsilon$ | (we )will loved | $\theta \alpha$ $\alpha \gamma \alpha \pi \eta$ Өоv́ $\mu \varepsilon$ | (we) will love | $\theta \alpha$ غ́ $\chi о ч \mu \varepsilon$ $\alpha \gamma \alpha \pi \eta \theta \varepsilon i ́$ | (we) will have been loved |
|  | $\theta \alpha$ $\alpha \gamma \alpha \pi$ ıó $\sigma \alpha \sigma \tau \varepsilon$ | (you)will be loved | $\theta \alpha$ $\alpha \gamma \alpha \pi \eta \theta \boldsymbol{\varepsilon i ́ t} \boldsymbol{\varepsilon}$ | (you) will love | $\theta \alpha \varepsilon ́ \chi \varepsilon \tau \varepsilon$ $\alpha \gamma \alpha \pi \eta \theta \varepsilon$ í | (you) will have been loved |
|  | $\theta \alpha \alpha \gamma \alpha \pi \mathbf{ı}$ ои́vtal | (they) will be loved | $\theta \alpha$ $\alpha \gamma \alpha \pi \eta \boldsymbol{\theta o v ́ v}$ | (they) will love | $\theta \alpha$ ह́ $\chi \circ v v$ $\alpha \gamma \alpha \pi \eta \theta \varepsilon i ́$ | (they) will have been loved |

## Appendix B: Evaluation of Modified Algorithm Testing Output

Number of words: 562242
Number of Stems: 138629
Frake's statistics:
Mean number of words per conflation class : 4.05573148475
Index compression factor : 0.75343535346
Unchanged Words : 7309 ------> ratio 0.0130720262642 (unchanged words /total words)

Mean Modified Humming Distance 2.44133602915
Median Modified Humming Distance: 2

| The number next to every word notes the modified Hamming distance according to Frakes from its stem |  |  |  |
| :---: | :---: | :---: | :---: |
| ---------ABA $\Theta------$ <br> is a stem for 12 words | ---------KППАNIZ------- | - --------ПРОП $\Omega \Lambda$------- | $\text { is a stem for } 50$ |
| ABA@ $\operatorname{sN} 2$ | KопANIZя 1 копanizoyn 3 | ПРОПএ^এNTA乏 5 ПРОП $\Omega \Omega \Omega 1$ | поү^תNTAE 5 |
| ABA@OYz 3 | копANIzoyme 4 | пРоПЛлOYees 5 | поүл 1 |
| ABA@OY 21 | копanizosoyn 5 | ПРоПЛлОүгЕ 4 | поүлOYeez 5 |
| ABA@OE 2 | Köanizosaite 6 | ПРОПДлОугаТЕ 6 | поулоугe 4 |
| ABA@OI 2 | KOIANIZONTAE 5 | ПРОПДлоүгAME 6 | поүnoyeate 6 |
|  | копANizontai 5 | ПРОПДлОүга 4 | поүмOYzame 6 |
| ABA@O 1 | KOпANIZOMAETE 6 | пропллоуntai 6 | поулоугa4 |
| ABA@Hट 2 | копAnizomat 4 | пРопллоуn 3 | пoynoyn 3 |
| ABA@H 1 | копanizete 3 | ПРоПллочME 4 | поулочме 4 |
| ABA@EL2 | Komanizetai 4 | ПРОПЛлОYMAгTE 7 | пoyaign 3 |
| ba@eli 3 | Koпanizezte 4 | ПРОПДлОYМАІ 5 | поyaioyntai 7 |
|  | Koпanizezai 4 | ПРОПДлнг 3 | поулioy 3 |
| Aba@E 1 | KOпANIzEE 2 | ПРОПДлНЕТЕ 4 | поуniontai 6 |
| Aba@A 1 | копANizeis 3 | ПРОПऽлНгО 4 | поүaietai 5 |
|  | копANIZEI 2 | ПРОПЛлНгЕТе 5 | поyaieste 5 |
| --АВАөоуллт------- | копanize 1 | ПРОПДлНЕЕ 3 | moyaiezai 5 |
| stem for 11 words | KOпANIZAN 2 | ПРОПЛлНटаTE 5 | поуaiez 3 |
|  | KOпANIZAME 3 | ПРОПД^HLAME 5 | поуaiemai 5 |


| ABA@OY | KОПANIZA 1 | ПРОПЛАНГА 3 | ПOYNIAE 3 |
| :---: | :---: | :---: | :---: |
|  |  | ПРОПЛ^НMENOᄃ 6 | ПОY^IA 2 |
| ABA@OY | ---------KOПANIZOT------- | ПРОПЛАНМЕNO 5 | ПOYAI 1 |
|  | is a stem for 1 words | ПРОПЛЛНМЕNН 5 | ПОҮЛНЕ 3 |
| АВА@ОY |  | ПРОПЛ^HMENE 5 | ПОҮАНГTE 4 |
| АВА@ОY | КОПАNIZOTAN 2 | ПРОПЛАНӨО 3 | ПОҮАНГOY 4 |
| ABA@OY $\Omega$ TOI 2 |  | ПРОПЛЛНӨНКЕГ 6 | ПОY^HГETE 5 |
| ABA@OY $\Omega$ RTO 1 | ---------KOПANIZOLALT------- | ПРОПЛ^НӨНКЕ 5 | ПОҮ^HटE 3 |
| ABA $\Theta О Y \wedge \Omega T H \Sigma 2$ | is a stem for 1 words | ПРОПЛАНӨНКА 5 | ПОҮ^НГATE 5 |
|  |  | ПРОПЛЛЕІТЕ 4 | ПОҮАНГAME 5 |
| ABA@OY $\Omega$ TH 1 | КОПANIZOEAETAN 2 | ПРОПЛ^ЕІГ 3 | ПОҮ^НГА 3 |
| ABA@OY |  | ПРОПИ^ЕІ 2 | ПОY^HMENOE 6 |
| ABA@OY | ---------KОПANIZONTOYГ------- |  | ПОҮЛНMENO 5 |
| ABA | is a stem for 1 words | --------ПРОПЛЛОҮТ------- | ПОYАНMENH 5 |
|  |  | is a stem for 1 words | ПОҮ^HMENE 5 |
|  | КОПANIZONTOYエAN 2 |  | $П О Ү \wedge Н \Theta \Omega ~ 3 ~$ |
| ---------АВАММО |  | ПРОПЛЛОҮTAN 2 | ПОҮАНӨНКЕГ 6 |
| is a stem for 11 words | ---------KOПANIZONT------- |  | ПОҮЛНӨНКЕ 5 |
|  | is a stem for 1 words | --------ПРОПЗЛОҮГ------- | ПОҮАНӨНКА 5 |
| АВАЄМО |  | is a stem for 2 words | ПОҮЛА 2 |
|  | КОПANIZONTAN 2 |  | ПОҮ ${ }^{\text {¢ }}$ ¢ 2 |
| АВАЄМОЛОГНТОҮГ 3 |  | ПРОПЛАОҮГOYN 3 | ПОY^ANE 3 |
| АВАӨМОЛОГНТОҮ 2 | ---------KОПANIZOM------- | $П Р О П \Omega \Lambda О Ү \Sigma A N ~ 2 ~$ | ПОҮ^AN 2 |
| АВАӨМО | is a stem for 1 words |  | ПОY ${ }^{\text {AME }} 3$ |
| АВАЄМОЛОГНТОІ 2 |  | --------- | ПОҮ^AEI 3 |
| АВАӨМОЛОГНТО 1 | КОПANIZOMOYN 3 | ПРОПЛЛОYГАГТ------- | ПОҮЛАГЕГ 4 |
| АВАЄМОЛОГНТНГ 2 |  | is a stem for 1 words | ПОҮЛАГЕ 3 |
|  | ---------KOПANIZOMAइT------- |  | ПОҮЛАГАТЕ 5 |
| АВАЄМОЛОГНТН 1 | is a stem for 1 words | ПРОПИЛОҮГАГTAN 2 | ПОҮЛАГАNЕ 5 |
| АВАӨМОЛОГНТЕГ 2 |  |  | ПОҮКАГАМЕ 5 |
| АВАӨМОЛОГНТЕ 1 | KOПANIZOMAइTAN 2 | ---------ПРОПЗ | ПОҮЛАГА 3 |
| АВАЄМОЛОГНТА 1 |  | is a stem for 1 words | ПОҮ^А 1 |
|  | ----------КОПANIZAT------- <br> is a stem for 1 words | ПРОПЛ^ОYNTAN 2 | ---- |
| ----------ABA@MI $\triangle \Omega$ T------- is a stem for 11 words | KOПANIZATE 1 | -- | ПОҮАХЕРІАР------- <br> is a stem for 1 words |
| is a stem for 11 words | ---------KOПANHइ------- | ПРОПЛАОҮМА <br> is a stem for 1 words | ПОҮАХЕРIAP 0 |
| ABA@MI $\triangle \Omega$ T $\Omega \mathrm{N} 2$ | is a stem for 6 words |  | IOYMXEPIAP 0 |
| ABA@Mİ 2 TOY $\Sigma 3$ |  | $П Р О П \Omega \wedge О Ү M A \Sigma T A N ~ 2 ~$ | ------- |
| ABA@MI $\triangle \Omega$ TOY 2 | КОПANHEOYN 3 |  | ПОҮАХЕР------- |
| ABA@MI $\triangle \Omega$ TOE 2 | KOПANHEOYME 4 | ---------ПРОПЛАНГ------- | is a stem for 1 words |
| ABA@MI $\triangle \Omega$ TOI 2 | KOПANHEEE 2 | is a stem for 10 words |  |
| ABA@MI $\triangle \Omega$ TO 1 | KOПANHLEİ 3 |  | ПОҮАХЕРIA 2 |
|  | КОПANHEEI 2 | ПРОПЛЛНГOYN 3 |  |
| ABA@MI $\triangle \Omega$ THE 2 | KOПANHEAN 2 | ПРОПЛАНГOYME 4 | -------- |
| ABA@MI $\Delta \Omega$ TH 1 |  |  | ПОҮЛОҮГ------- |
| ABA $\Theta$ MI $\triangle \Omega$ TE $\Sigma 2$ |  |  | is a stem for 1 words |
| ABA@MI $\triangle \Omega$ TE 1 |  | ПРОПЛАНГЕ $\Omega 3$ |  |
| ABA@MI $\triangle \Omega$ TA 1 |  | ПРОПЗЛНГЕ | ПОҮ^OYГAN 2 |



The previous table is a small sample from an extensive list of conflation classes created by the modified stemmer. In this small list we included examples of both understemming and overstemming. One overstemming example is the conflation class of "ПOY "" which includes words from the word " $\pi$ оv $\lambda i$ " (bird) and " $\pi$ оv $\lambda \dot{\alpha} \omega$ " (to sell). On the contrary, understemming examples have to do with the grammatical changes that words undergo in past tenses.

## Appendix B: On-line Stemmer



Our implementation of the algorithm is freely available at http://gelaligo.org/stemmer under an LGPL licence, along with an on-line demo.

## Appendix C: Stop Word List

| ABA | ANAMEEA | AYPIO |
| :---: | :---: | :---: |
| АГА | ANAMETAEY | АФН |
| АГН | ANEY | АФОТОУ |
| $A \Gamma \Omega$ | ANTI | АФОУ |
| A $\triangle \mathrm{H}$ | ANTIIEPA | AX |
| $\mathrm{A} \Delta \Omega$ | ANTIL | AXE |
| AE | ANS | AXO |
| AEI | AN $\Omega$ TEP $\Omega$ | A $\Psi \mathrm{A}$ |
| $\mathrm{A} \Theta \Omega$ | A | A ${ }^{\text {PE }}$ |
| AI | АП | A $\Psi \mathrm{H}$ |
| AIK | АПЕNANTI | A $\Psi Y$ |
| AKH | АПО | $\mathrm{A} \Omega \mathrm{E}$ |
| AKOMA | АПОЧЕ | A $\Omega \mathrm{O}$ |
| AKOMH | АП $\Omega$ | BAN |
| AKPIB $\Omega \Sigma$ | APA | BAT |
| A $\wedge$ A | АРАГЕ | BAX |
| A | APE | BEA |
| A | APK | BEBAIOTATA |
| A $\triangle \Lambda A X O Y$ | APKETA | BHE |
| $\mathrm{A} \Lambda \Lambda I \Omega \Sigma$ | AP」 | BIA |
| A $\wedge \wedge$ ISTIKA | APM | BIE |
| A $\Lambda \Lambda O I \Omega \Sigma$ | APT | BIH |
| A $\Lambda \Lambda O I \Omega T I K A$ | APY | BIO |
| A $\Lambda \Lambda O T E$ | $\mathrm{AP} \Omega$ | BOH |
| A $\wedge$ T | A $\Sigma$ | BO $\Omega$ |
| $\mathrm{A} \Lambda \Omega$ | A $\Sigma \mathrm{A}$ | BPE |
| AMA | A SO | ГА |
| AME | ATA | $Г А В$ |
| AMESA | ATE | ГАР |
| AME $2 \Omega \Sigma$ | ATH | ГEN |
| AM $\Omega$ | ATI | ГEL |
| AN | ATM | ГН |
| ANA | ATO | ГHN |


| ГI | $\Delta Y E$ | E $\Xi$ |
| :---: | :---: | :---: |
| ГІА | $\Delta \mathrm{YO}$ | E®AФNA |
| ГIE | $\Delta \Omega$ | E $\Xi \mathrm{I}$ |
| $\Gamma \mathrm{IN}$ | EAM | EEILOY |
| ГІО | EAN | $\mathrm{E} \Xi \Omega$ |
| ГКI | EAP | EOK |
| ГКY | E@H | ЕПAN $\Omega$ |
| ГОН | EI | ЕПЕİH |
| ГОО | EIDEMH | ЕПЕІТА |
| ГРНГОРА | EI@E | ЕПН |
| ГРІ | EIMAI | ЕПI |
| ГРY | EIMALTE | ЕПİНГ |
| ГYН | EINAI | ЕПОМЕNתГ |
| $Г \mathrm{YP} \Omega$ | EIL | EPA |
| $\Delta \mathrm{A}$ | EILAI | E $\Sigma$ |
| $\Delta \mathrm{E}$ | EİAETE | ELAE |
| $\Delta \mathrm{EH}$ | EILTE | ELE |
| $\Delta \mathrm{EI}$ | EITE | EऽEIL |
| $\Delta \mathrm{EN}$ | EIXA | ESENA |
| $\Delta \mathrm{E} \Sigma$ | EIXAME | E 2 H |
| $\Delta \mathrm{H}$ | EIXAN | E $2 T \Omega$ |
| $\Delta \mathrm{H} \Theta \mathrm{EN}$ | EIXATE | E 5 Y |
| $\Delta \mathrm{H} \Lambda \mathrm{A} \triangle \mathrm{H}$ | EIXE | $E \Sigma \Omega$ |
| $\Delta \mathrm{H} \Omega$ | EIXE | ETI |
| $\Delta \mathrm{I}$ | EK | ETEI |
| $\triangle \mathrm{IA}$ | EKO | EY |
| $\Delta \mathrm{IAPK} \Omega \Sigma$ | EKEI | EYA |
| $\triangle \mathrm{IO}$ 人OY | E $\Lambda$ A | ЕҮГЕ |
| $\Delta I \Sigma$ | E $\Lambda$ I | EY@Yг |
| $\Delta I X \Omega \Sigma$ | ЕМП | EYTYX $\Omega$ |
| $\Delta \mathrm{O} \Lambda$ | EN | EФE |
| $\triangle \mathrm{ON}$ | ENTE $\Lambda \Omega \Sigma$ | ЕФE®HГ |
| $\triangle \mathrm{PA}$ | ENTOE | ЕФТ |
| $\Delta \mathrm{PY}$ | ENTQMETAEY | EXE |
| $\Delta \mathrm{PX}$ | ENת | EXEI |


| EXEIL | IBO | КАП®इ |
| :---: | :---: | :---: |
| EXETE | $\mathrm{I} \Delta \mathrm{H}$ | KAT |
| EX@Eร | $\mathrm{I} \triangle \mathrm{I} \Omega \Sigma$ | KATA |
| EXOME | IE | KATI |
| EXOYME | II | KATITI |
| EXOYN | III | KATOMIN |
| EXTE ${ }^{\text {c }}$ | IKA | KAT $\Omega$ |
| EX $\Omega$ | $\mathrm{I} \Lambda \mathrm{O}$ | KA $\Omega$ |
| E $\Omega \Sigma$ | IMA | KBO |
| ZEA | INA | KEA |
| ZEH | $\mathrm{IN} \Omega$ | KEI |
| ZEI | $\mathrm{I} \Xi \mathrm{E}$ | KEN |
| ZEN | $\mathrm{I} \Xi \mathrm{O}$ | KI |
| ZHN | IO | KIM |
| Z $\Omega$ | IOI | KIO^A |
| H | İA | KIT |
| $\mathrm{H} \Delta \mathrm{H}$ | ILAME | KIX |
| $\mathrm{H} \Delta \mathrm{Y}$ | ILE | KKE |
| $\mathrm{H} \Theta \mathrm{H}$ | I I H | K $\wedge$ ILE |
| Н $\Lambda \mathrm{O}$ | ILIA | К $\wedge \Pi$ |
| HMI | ILO | KOK |
| НПА | I $\Sigma \Omega \Sigma$ | KONTA |
| HzAETE | I I B | KOX |
| HEOYN | $\mathrm{I} \Omega \mathrm{N}$ | KT^ |
| HTA | $I \Omega \Sigma$ | KYP |
| HTAN | $\mathrm{I} \alpha \nu$ | KYPI $\Omega \Sigma$ |
| HTANE | KA $\Theta$ | $\mathrm{K} \Omega$ |
| HTOI | KA@E | $\mathrm{K} \Omega \mathrm{N}$ |
| HTTON | KA@ETI | $\Lambda \mathrm{A}$ |
| $\mathrm{H} \Omega$ | KA@O^OY | $\Lambda$ EA |
| $\Theta \mathrm{A}$ | KA@ ${ }^{\text {I }}$ | $\Lambda \mathrm{EN}$ |
| OYE | KAI | \EO |
| $\Theta \Omega P$ | KAN | NIA |
| I | КАПОТЕ | КІГАКІ |
| IA | КАПОҮ | АІГО |


| $\Lambda$ ІГ $\Omega$ TEPO | MHN | NTA |
| :---: | :---: | :---: |
| 人IO | МНП $\Omega \Sigma$ | NTE |
| NIP | MHTE | NTI |
| $\Lambda О Г \Omega$ | MI | NTO |
| ЛОІПА | MIE | NYN |
| $\Lambda$ OIMON | MIL | $\mathrm{N} \Omega \mathrm{E}$ |
| 人OE | MME | N 2 PIL |
| $\Lambda \Sigma$ | MNA | EANA |
| $\Lambda Y \Omega$ | MOB | ЕAФNIKA |
| MA | MO^IL | $\Xi E \Omega$ |
| MAZI | MOAONOTI | $\Xi I$ |
| MAKAPI | MONAXA | O |
| MA^İTA | MONOMIAE | OA |
| MA $\Lambda \Lambda O N$ | MOY | ОАП |
| MAN | МПА | $\mathrm{O} \Delta \mathrm{O}$ |
| MAE | МПОРЕІ | OE |
| MAE | M | OZO |
| MAT | МПРАВО | OHE |
| ME | МПРОГ | OI |
| ME@AYPIO | МП $\Omega$ | OIA |
| MEI | MY | OIH |
| MEION | MYA | OKA |
| ME $\Lambda$ | MYN | ОЛОГҮРА |
| ME^EI | NA | OAONEN |
| ME $\Lambda \Lambda$ ETAI | NAE | O |
| MEMIA | NAI | $\mathrm{O} \Lambda \Omega \Sigma \Delta \mathrm{IO} \Lambda \mathrm{OY}$ |
| MEN | NAO | OM $\Omega \Sigma$ |
| MES | $\mathrm{N} \Delta$ | ON |
| MESA | NE ${ }^{\circ}$ | ONE |
| MET | NI | ONO |
| META | NIA | ОПА |
| METAEY | NIK | ОПЕ |
| MEXPI | NIN | ОПН |
| MH | NIN | ОПО |
| $\mathrm{MH} \Delta \mathrm{E}$ | NIO | OПОІАДНПОТЕ |


| OПOIAN $\triangle$ HПОTE | OTE | ПІК |
| :---: | :---: | :---: |
| ОПОІАऽДНПОТЕ | OTI | $\Pi$ ПО |
| OПОІАНПОТЕ | OTIДНПОТЕ | $\Pi I \Sigma \Omega$ |
| OПOIE $\triangle$ НПОТЕ | OY | ПIT |
| ОПОІОДНПОТЕ | OYAE | $\Pi I \Omega$ |
| OПОІОNДНПОТЕ | OYK | ПИAI |
| OПOIOг НПОТЕ $^{\text {a }}$ | OY̌ | П^EON |
| ОПОІОУДНПОТЕ | OYTE | ПАНN |
| ОПОІОҮГДНПОТЕ | OYФ | $\Pi \Lambda \Omega$ |
| ОПОІЛNДНПОТЕ | OXI | ПМ |
| ОПОТЕДНПОТЕ | ОЧА | ПОА |
| ОПОУ | OЧE | ПОЕ |
| ОПОҮДНПОТЕ | ОЧН | ПОД |
| ОП $\Omega \Sigma$ | OЧI | ПОАУ |
| OPA | ОЧО | ПОП |
| OPE | ПА | ПОТЕ |
| OPH | ПААI | ПОҮ |
| OPO | ПАN | ПОҮ®Е |
| ОРФ | ПANTOTE | ПОY@ENA |
| OP $\Omega$ | ПANTOY | ПРЕПЕІ |
| OEA | ПАNT®इ | ПРІ |
| OऽA $\triangle$ НПОТЕ | ПАП | ПPIN |
| OLE | ПАР | ПРО |
|  | ПАРА | ПPOKEIMENOY |
| ОГНДНПОТЕ | ПЕІ | ПРОКЕITAI |
| OгHNДНПОTE | ПЕР | ПРОПЕРГІ |
| OこН $\triangle \triangle$ НПОТЕ | ПЕРА | ПРОГ |
| ОгОДНПОТЕ | ПЕРI | ПРОТОУ |
| OгOIДHПOTE | ПЕРІПОҮ | ПРОХ@ЕГ |
| OГONДНПОТЕ | ПЕРГI | ПРОХТЕ |
|  | ПЕРYГI | ПРОТҮТЕРА |
| ОГОYДНПОТЕ | ПЕг | ПҮА |
| OLOY $\triangle$ НПОТЕ | ПI | ПYЕ |
| Oг $2 \mathrm{~N} \triangle$ НПОТЕ | ПІА | ПҮО |
| OTAN | ПI@ANON | ПҮР |


| ПХ | SEP | TA |
| :---: | :---: | :---: |
| $\Pi \Omega$ | SET | TA $\Delta \mathrm{E}$ |
| $\Pi \Omega \Lambda$ | $\Sigma$ EФ | TAK |
| $\Pi \Omega \Sigma$ | ऽHMEPA | TAN |
| PA | $\Sigma \mathrm{I}$ | TAO |
| PAI | ᄃIA | TAY |
| РАП | $\Sigma$ ІГА | TAXA |
| PAE | $\Sigma \mathrm{IK}$ | TAXATE |
| PE | SIX | TE |
| PEA | $\Sigma \mathrm{KI}$ | TEI |
| PEE | гOI | TE $\Lambda$ |
| PEI | ऽOK | TEAIKA |
| PHE | ऽO^ | TEへIK $\Omega \Sigma$ |
| $\mathrm{P} \Theta \Omega$ | 5ON | TE $\Sigma$ |
| PIO | ऽOE | TET |
| PO | ऽOY | TZO |
| PO'̈̈ | $\Sigma \mathrm{PI}$ | TH |
| POE | ऽTA | TH^ |
| POZ | ऽTH | THN |
| POH | $\Sigma$ THN | THE |
| $\mathrm{PO} \Theta$ | ऽTHE | TI |
| POI | इTIL | TIK |
| POK | ETO | TIM |
| $\mathrm{PO} \Lambda$ | $\Sigma$ TON | ТІПОТА |
| PON | $\Sigma$ TOY | ТІПОТЕ |
| POE | इTOY | Tİ |
| POY | $\Sigma \mathrm{T}$ / N | TNT |
| EAI | $\Sigma \mathrm{Y}$ | TO |
| $\Sigma \mathrm{AN}$ | ГҮГХРОN $2 \Sigma$ | TOI |
| $\Sigma \mathrm{AO}$ | IYN | TOK |
| LA | IYNAMA | TOM |
| LE | $\Sigma$ YNEП $\Omega \Sigma$ | TON |
| SEIL |  | ТОП |
| EEK | $\Sigma \mathrm{XE}$ OON | TOE |
| EEE | $\Sigma \Omega \Sigma \mathrm{TA}$ | TOESN |

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| TOEA | ҮПОЧІN | XOH |
| :---: | :---: | :---: |
| TOEE | Y 2 TEPA | XO^ |
| TOEH | YФН | $\mathrm{XP} \Omega$ |
| TOLHN | YчН | XTE |
| TOEHE | ФА | X $\Omega$ PIL |
| TOEO | ФАї | X |
| TOEOI | ФАE | YE |
| TOEON | ФAN | $\Psi \mathrm{H} \Lambda \mathrm{A}$ |
| TOEOE | ФА $\Xi$ | $\Psi \mathrm{I}$ |
| TOEOY | ФAट | YIT |
| TOEOY | ФА $\Omega$ | $\Omega$ |
| TOTE | ФEZ | $\Omega \mathrm{A}$ |
| TOY | ФEI | $\Omega A \Sigma$ |
| TOY^AXIETO | ФETOE | $\Omega \Delta \mathrm{E}$ |
| TOY^AXIETON | ФEY | $\Omega E \Sigma$ |
| TOY̌ | ФI | $\Omega \Theta \Omega$ |
| T $\Sigma$ | ФI^ | $\Omega \mathrm{MA}$ |
| T $\Sigma \mathrm{A}$ | Фİ | $\Omega \mathrm{ME}$ |
| TEE | ФOE | $\Omega \mathrm{N}$ |
| TYXON | ФПА | $\Omega \mathrm{O}$ |
| $\mathrm{T} \Omega$ | ФРІ | תON |
| $\mathrm{T} \Omega \mathrm{N}$ | XA | $\Omega \mathrm{OY}$ |
| T $\Omega$ PA | XAH | $\Omega \Sigma$ |
| YAE | XA^ | $\Omega \Sigma A N$ |
| YBA | XAN | $\Omega \Sigma \mathrm{H}$ |
| YBO | ХАФ | תГOTOY |
| YIE | XE | $\Omega \Sigma \Pi О Ү$ |
| YIO | XEI | $\Omega \Sigma T E$ |
| $\mathrm{Y} \Lambda \mathrm{A}$ | X $\Theta E \Sigma$ | תГTOLO |
| Y $\$ H & XI & $\Omega$, |  |  |
| YNI | XIA | $\Omega X$ |
| ҮП | XIS | $\Omega \Omega \mathrm{N}$ |
| ҮПЕР | XIO | ГIATI |
| ҮПО | X $\wedge$ M |  |
| ҮПОЧН | XM |  |

